Chapter 2 Lectures

Stallings - 9th Ed

Applic	Application
Lib	programming interface Application
Ope	binary interface Instruction Set
Execu	Architecture
System	
I/O devices and networking	
System I/O devices and	

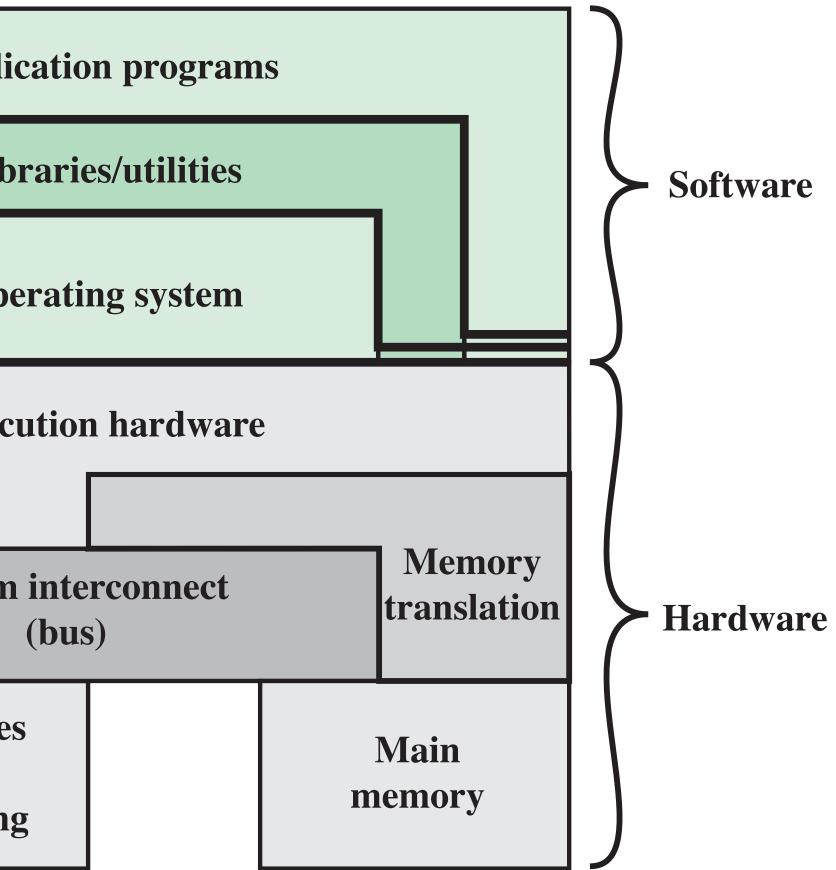


Figure 2.1 Computer Hardware and Software Structure

Operating System

application programs hardware

- A program that controls the execution of
- An interface between applications and

Operating System Objectives

- Convenience
- Efficiency
 - Allows computer system resources to be used in an efficient manner
- Ability to evolve
 - introduction of new system functions without interfering with service

– Makes the computer more convenient to use

– Permit effective development, testing, and

Services Provided by the Operating System

- Program development
 Editors and debuggers
- Program execution
- Access to I/O devices
- Controlled access to files
- Memory management
- System access
- Network support

Services Provided by the Operating System (2)

- Error detection and response
 - Internal and external hardware errors
 - Memory error
 - Device failure
 - Software errors
 - Arithmetic overflow
 - Access forbidden memory locations
 - Operating system cannot grant request of application

Services Provided by the Operating System (3)

- Accounting
 - Collect usage statistics
 - Monitor performance
 - Used to anticipate future enhancements
 - Used for billing purposes
 - check out the log files of a unix system
 - where do you find this information?

Operating System

programs) that is executed

Sequence 2

• Responsible for managing resources • An OS is just a program (or set of

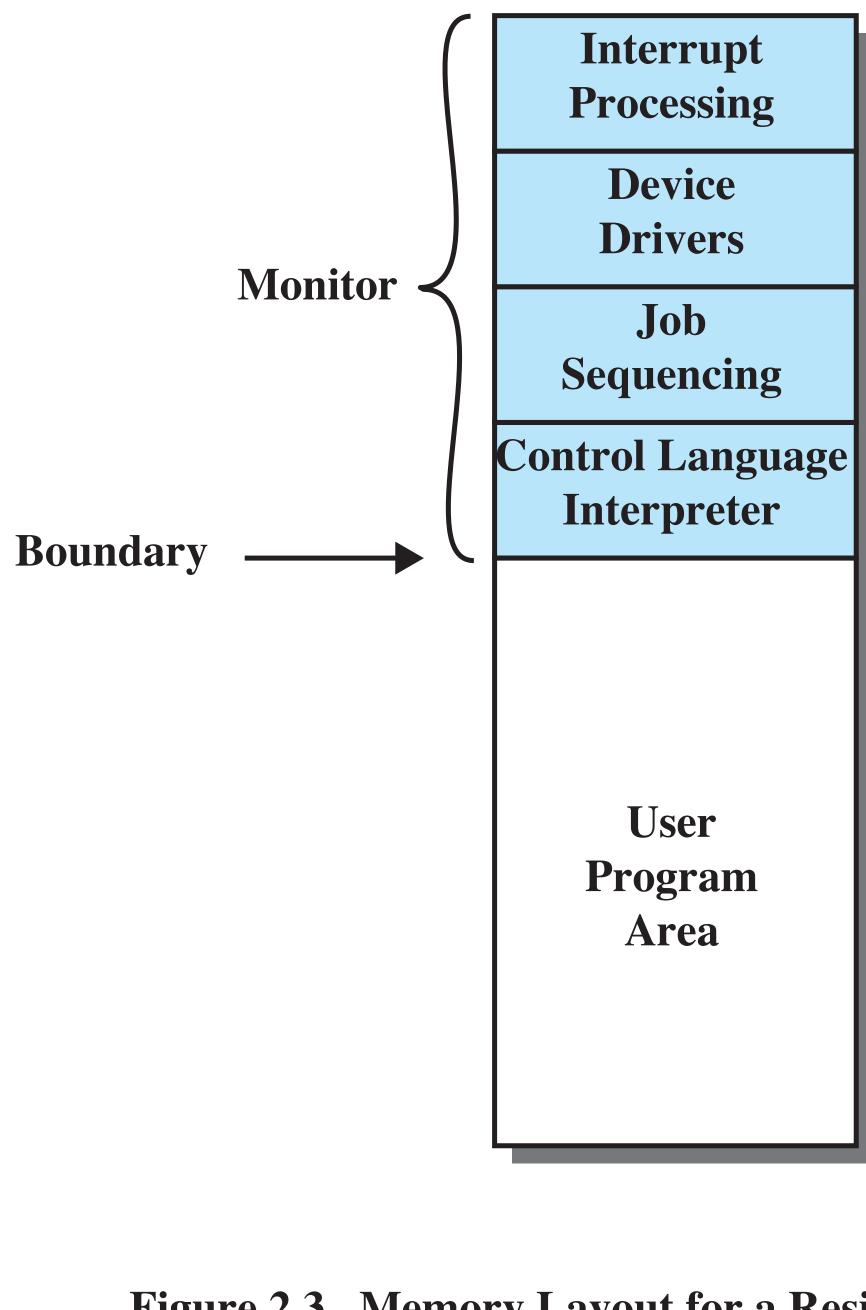


Figure 2.3 Memory Layout for a Resident Monitor

Evolution of Operating Systems

- Serial Processing
 - -No operating system
 - Machines run from a console with display lights, toggle switches, input device, and printer
 - Schedule time, e.g. sign up
 - Setup included loading the compiler and source program, saving compiled program, loading and linking

Evolution of Operating Systems • Simple Batch Systems

- Monitor

 - Software that controls the sequence of events • Batch jobs together
 - Program branches back to monitor when finished

Job Control Language (JCL)

- Special type of programming language • Provides instruction to the monitor, e.g. – What compiler to use
- - What data to use

Sequence 3

Hardware Features of Interest in Operating Systems

- User/Supervisor Mode
- Interrupts/Traps/Exceptions
- Real-time Clock (Timer) Operation
- Support for Virtual Memory and Memory Protection

Hardware Features

- Memory protection

 Do not allow the memory area containing the monitor to be altered

 Timer
 - Prevents a job from monopolizing the system

Hardware Features

Privileged instructions

Certain machine level instructions can only be executed by the monitor

Interrupts

Early computer models did not have this capability

Memory Protection

- Certain instructions may not be executed
- User program executes in **user mode** • Monitor executes in system mode
 - Kernel mode
 - Privileged instructions are executed - Protected areas of memory may be accessed

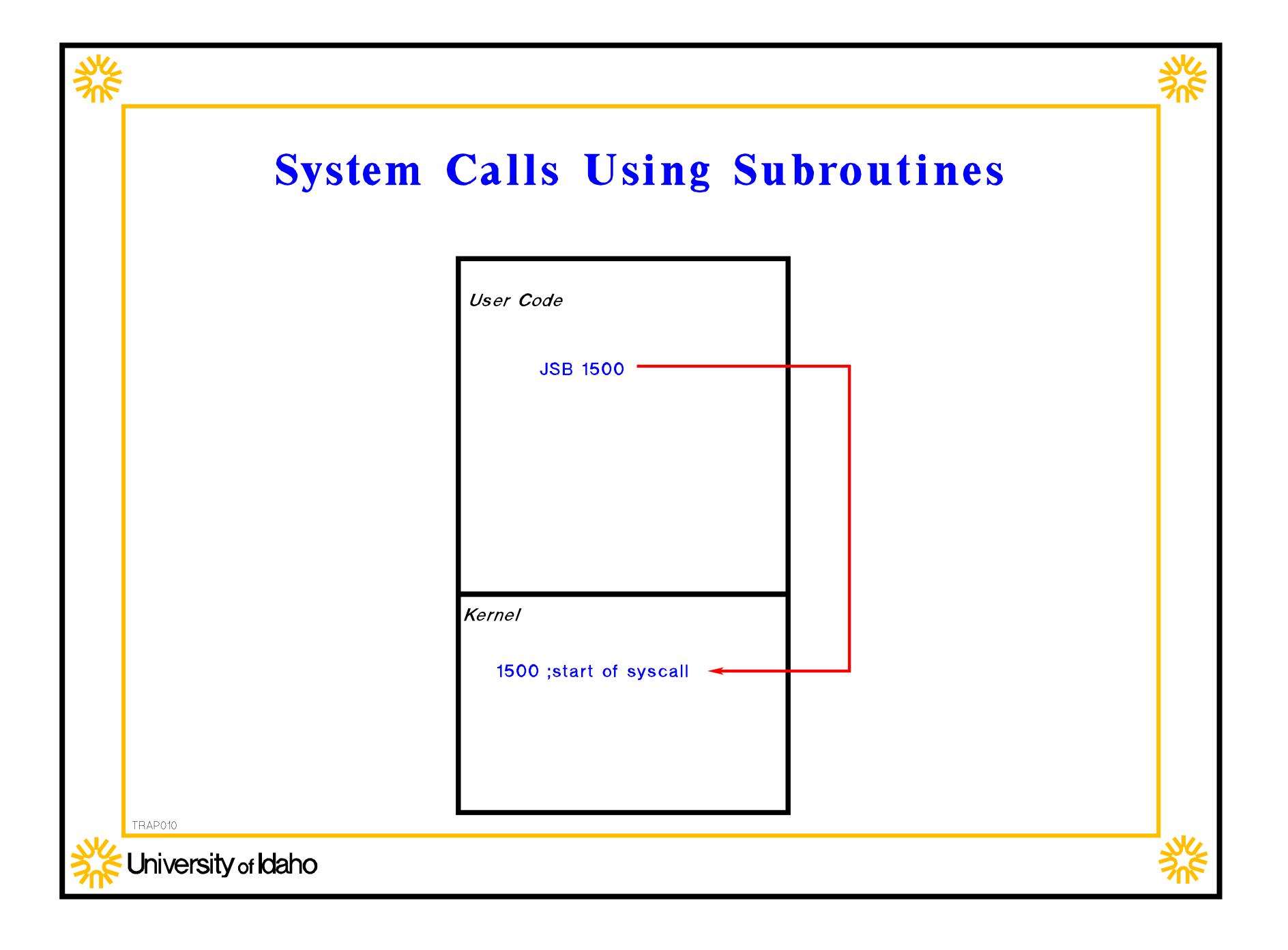
Traps and Interrupts and Exceptions (Oh My!)

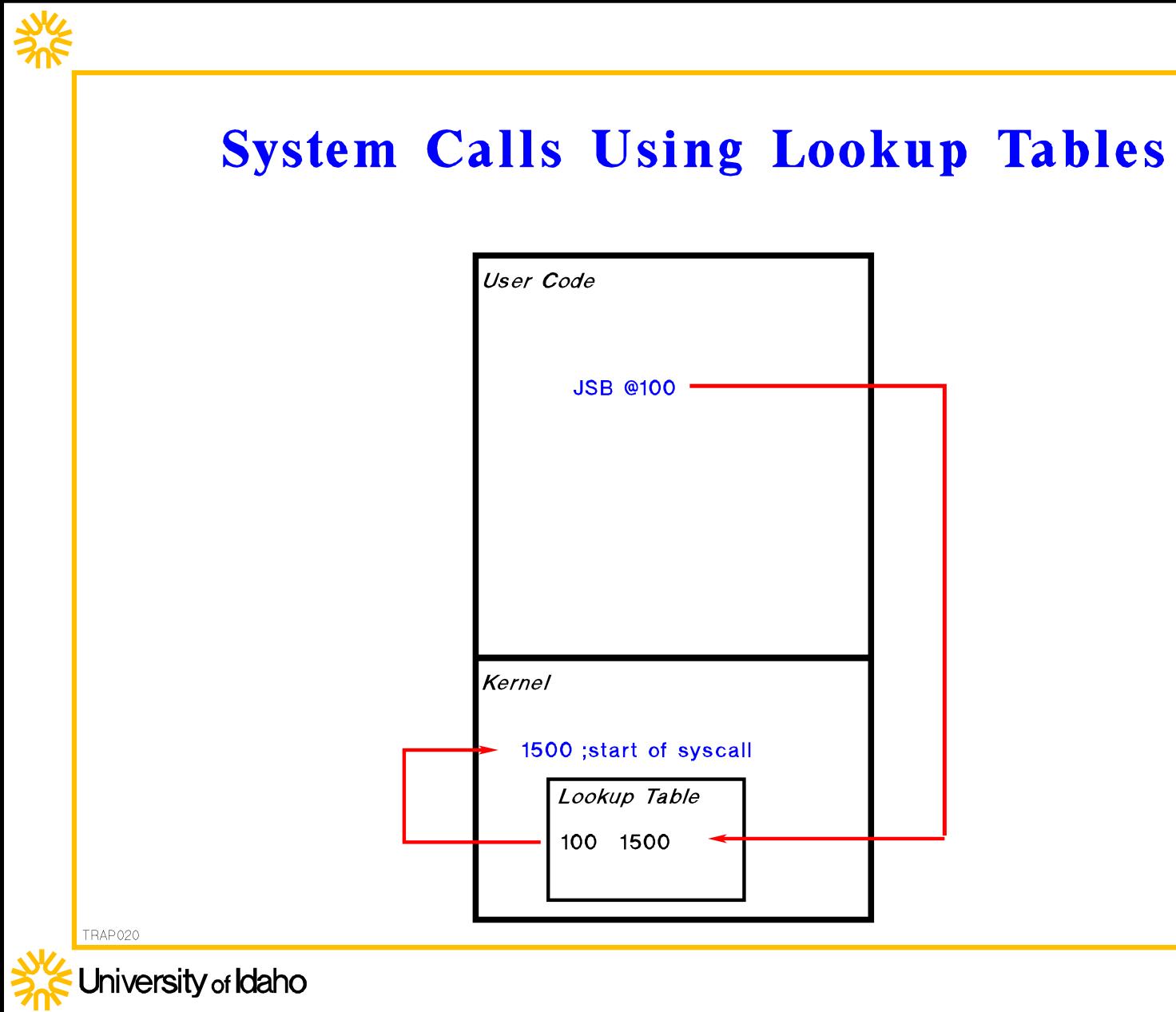
- <u>Trap</u> Caused by executing a "trap" instruction. Used to request an operation from the operating system. Also called a "syscall" or "synchronous trap."
- Interrupt caused by a hardware device when it needs service. Also called an "asynchronous trap."
- <u>Exception</u> caused when an illegal operation (eg, divide by zero or memory) reference out of bounds) is attempted.

subroutine call" through the vector table, along with a change from user to kernel mode.

With all three, the processor response is the same - an "indirect

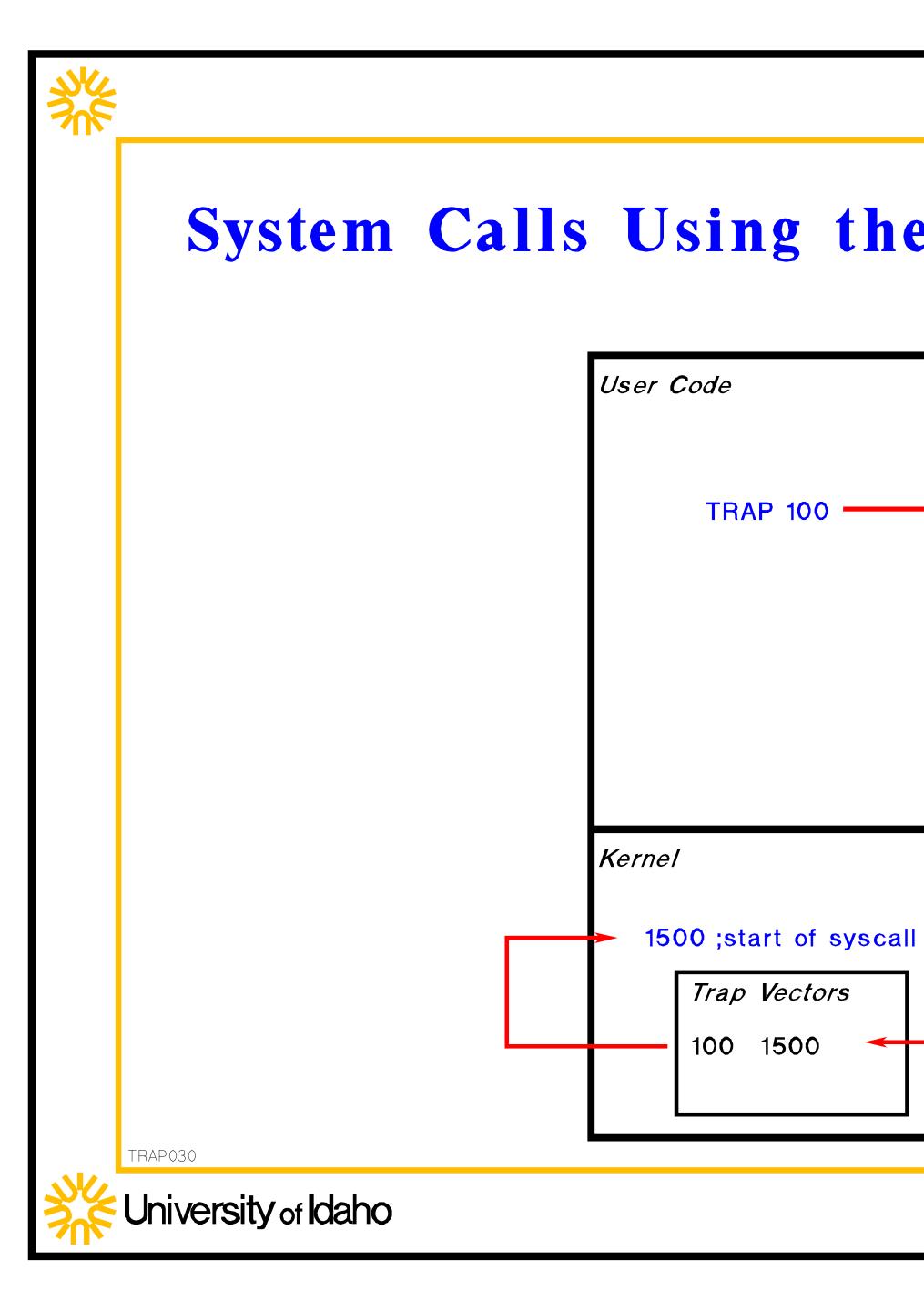


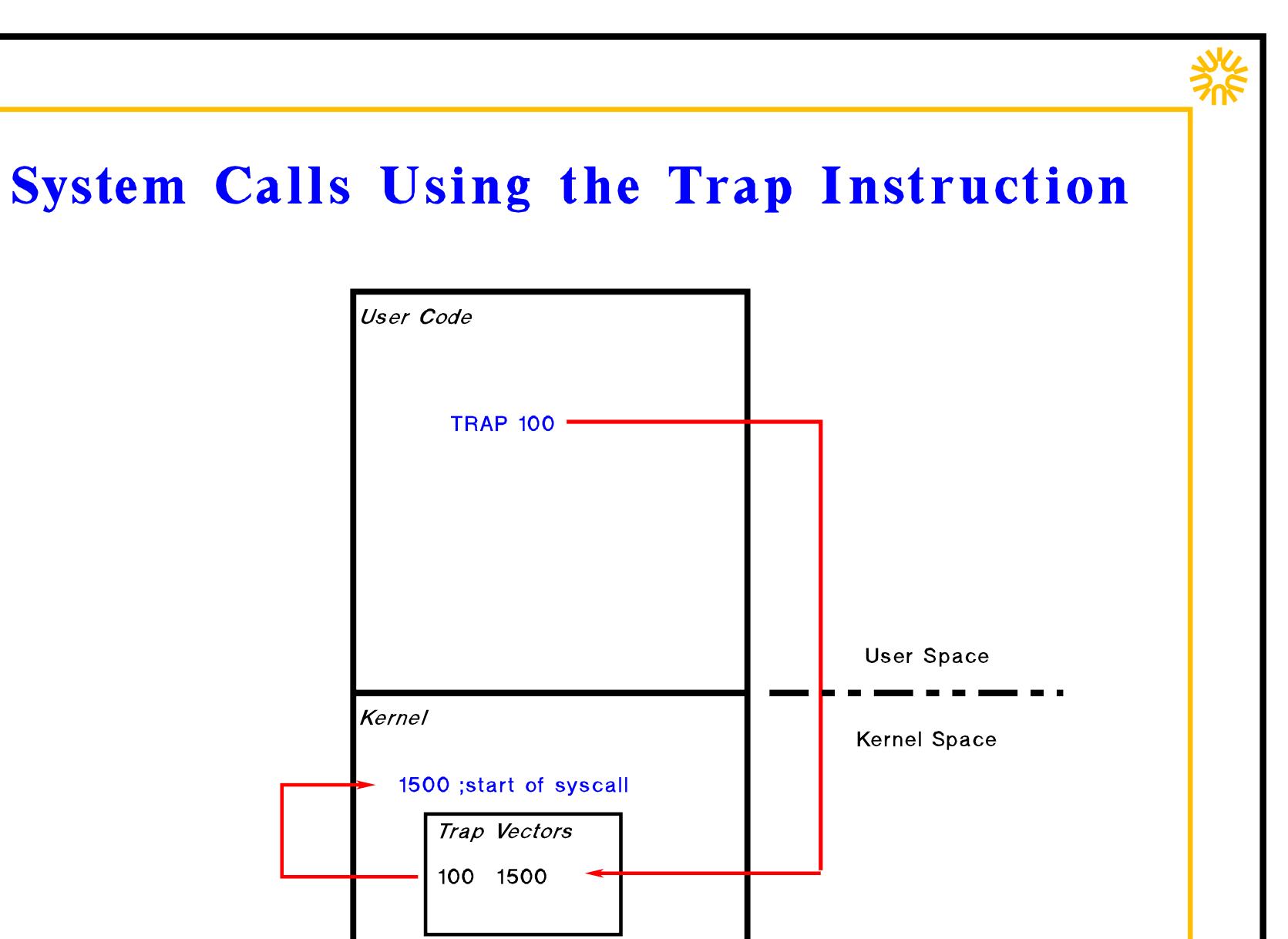












Read one record from Execute 100 instruction Write one record to file TOTAL

Percent CPU Utilizatio

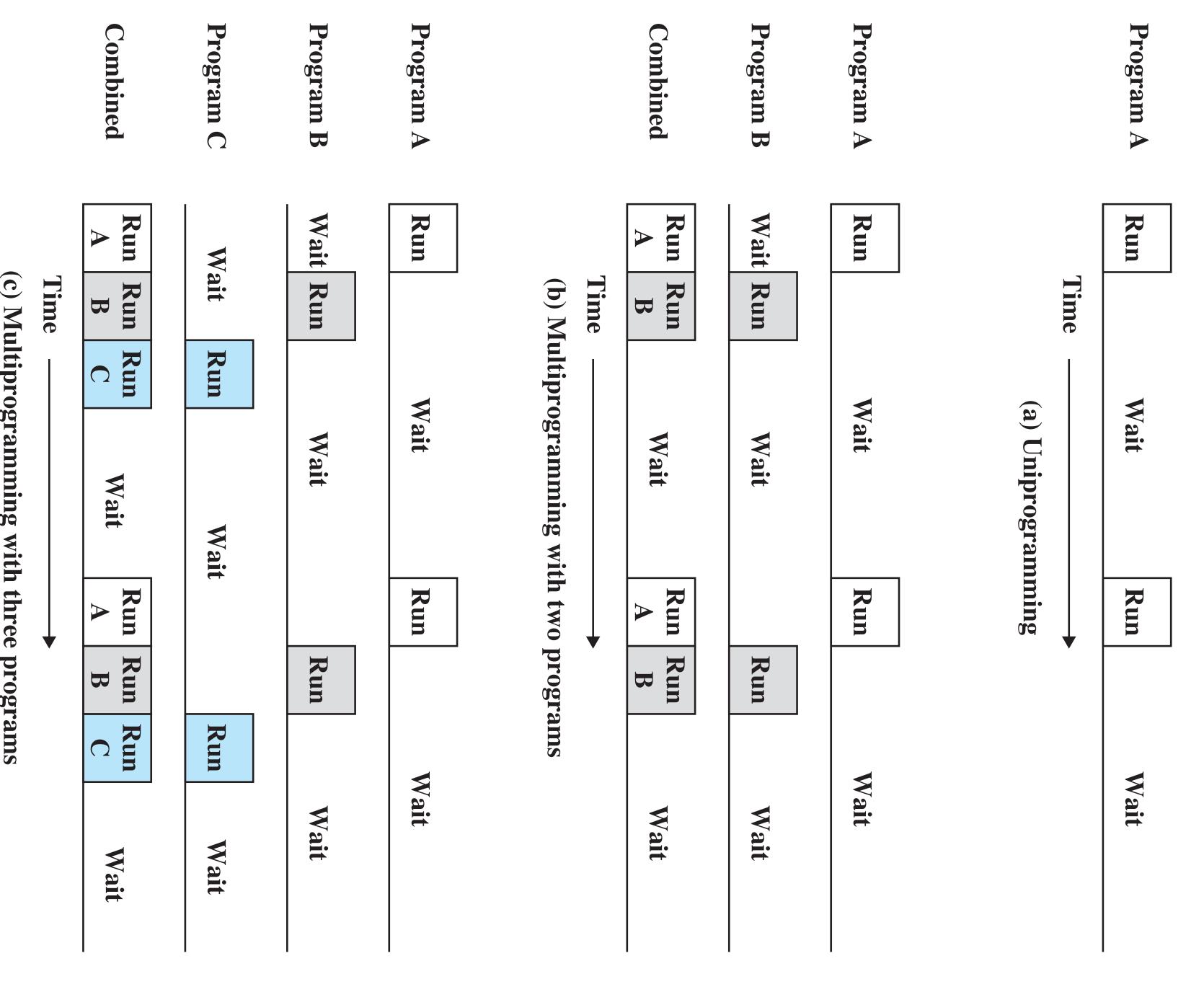
file
$$15 \ \mu s$$

ns $1 \ \mu s$
e $\frac{15 \ \mu s}{31 \ \mu s}$
on $= \frac{1}{31} = 0.032 = 3.2\%$

Figure 2.4 System Utilization Example

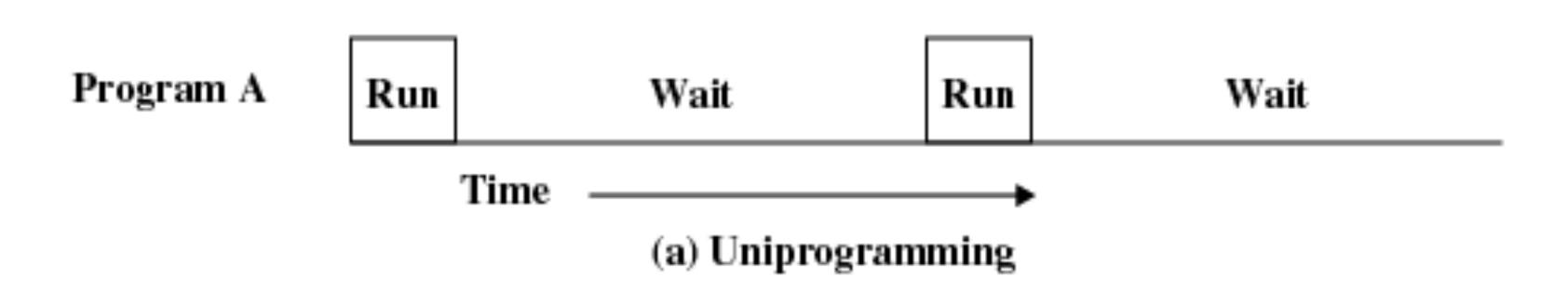
	Combined	Program C	Program B
	Run A	Wait	Wait Run
Time (c) Mu	Run Run B C	ait	Run
 ıltipro	Run C	Run	
Time	Wait Run A	Wait	Wait
e prog	Run Run Run A B C		Run
ams	Run C	Run	
	Wait	Wait	Wait

Figure 2.5 **Multiprogramming Example**



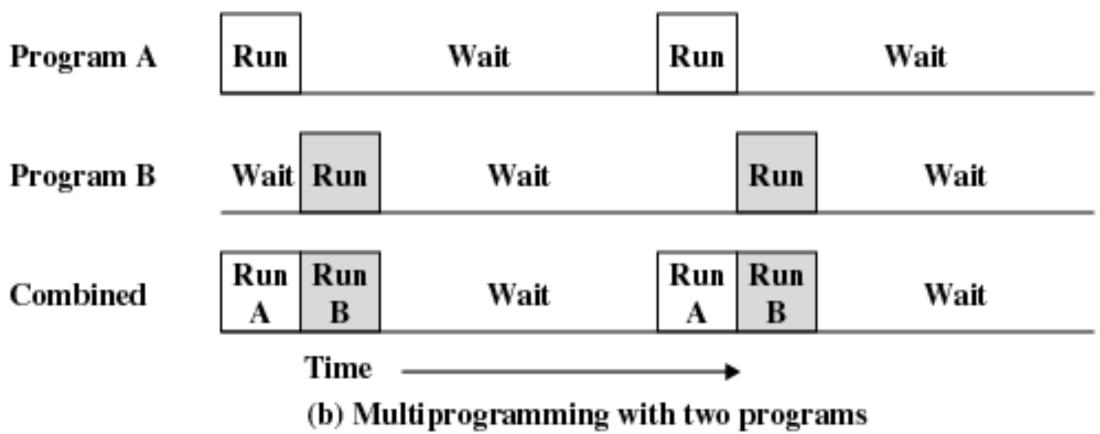
Uniprogramming

Processor must wait for I/O instruction to complete before preceding

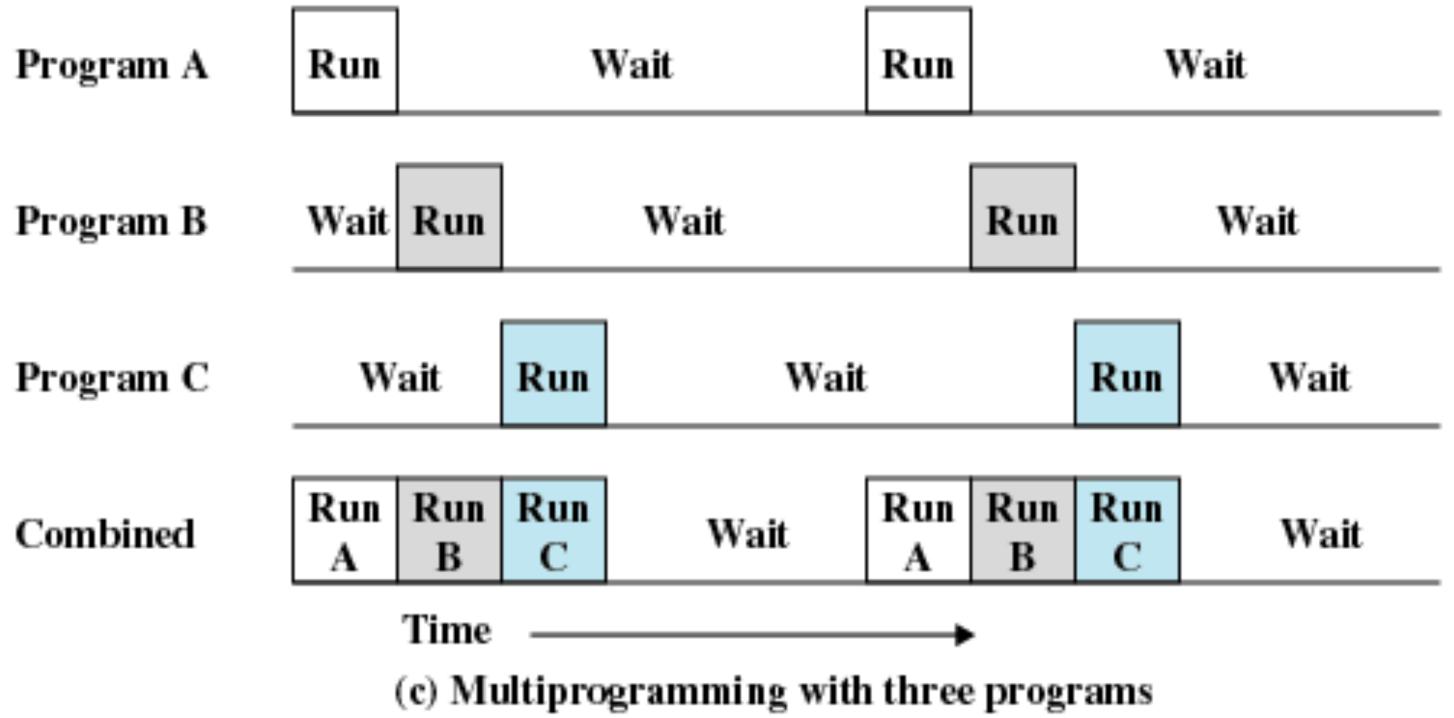


Multiprogramming

• When one job needs to wait for I/O, the processor can switch to the other job

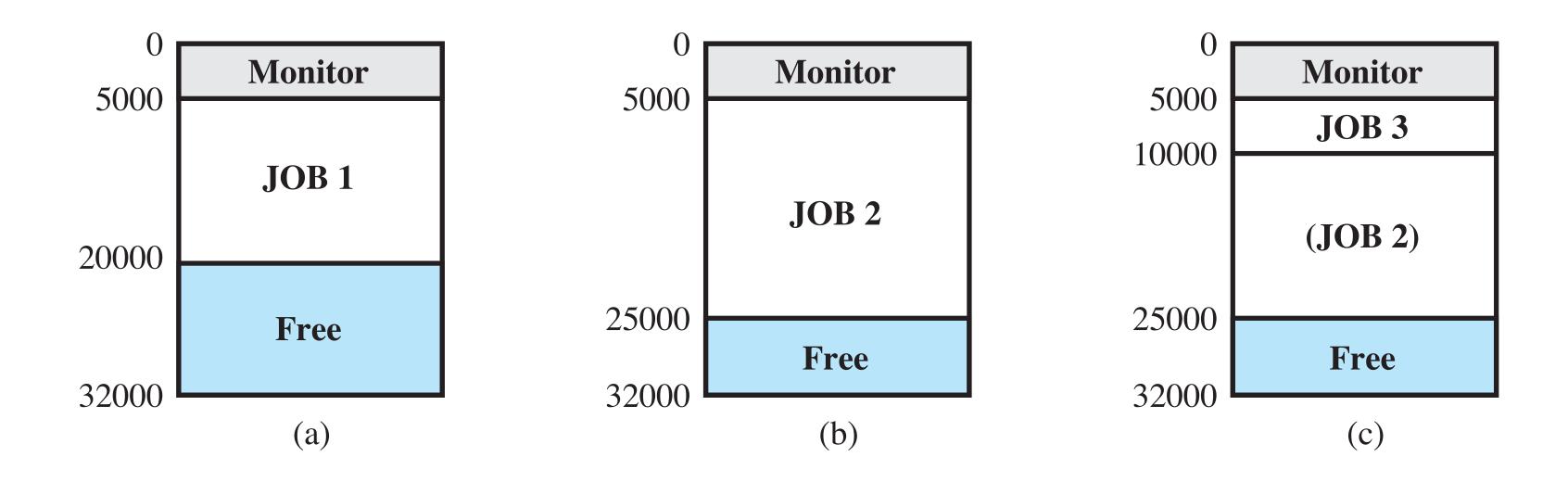


Multiprogramming



Time Sharing

- Using multiprogramming to handle multiple interactive jobs
- Processor's time is shared among multiple users
- Multiple users simultaneously access the system through terminals



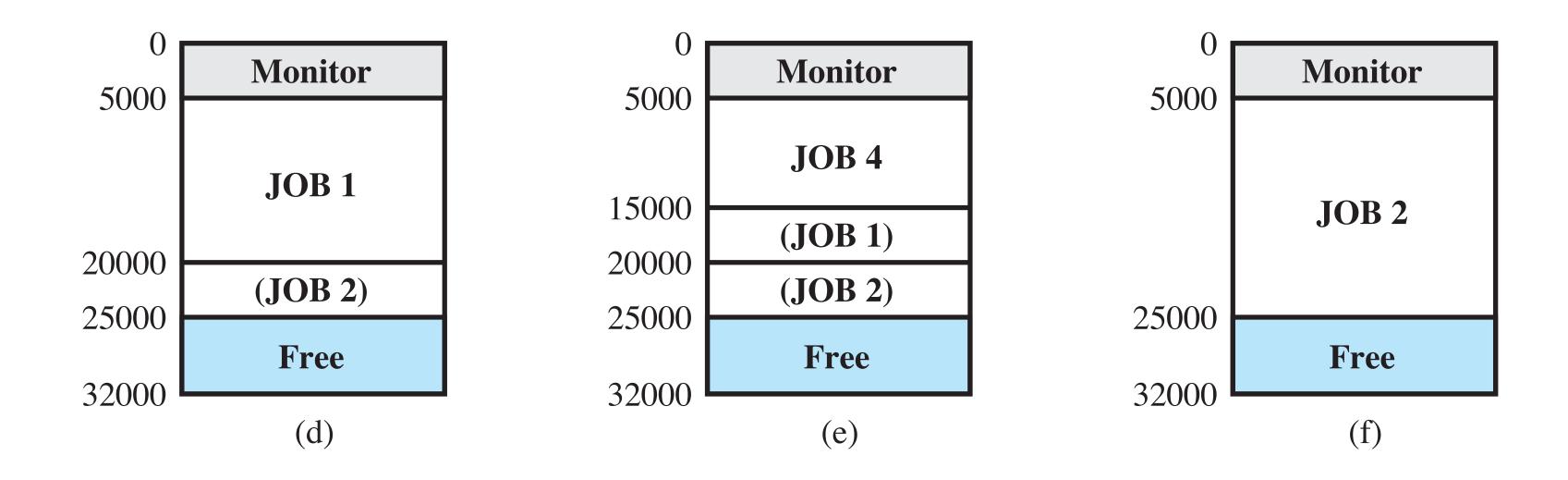


Figure 2.7 CTSS Operation

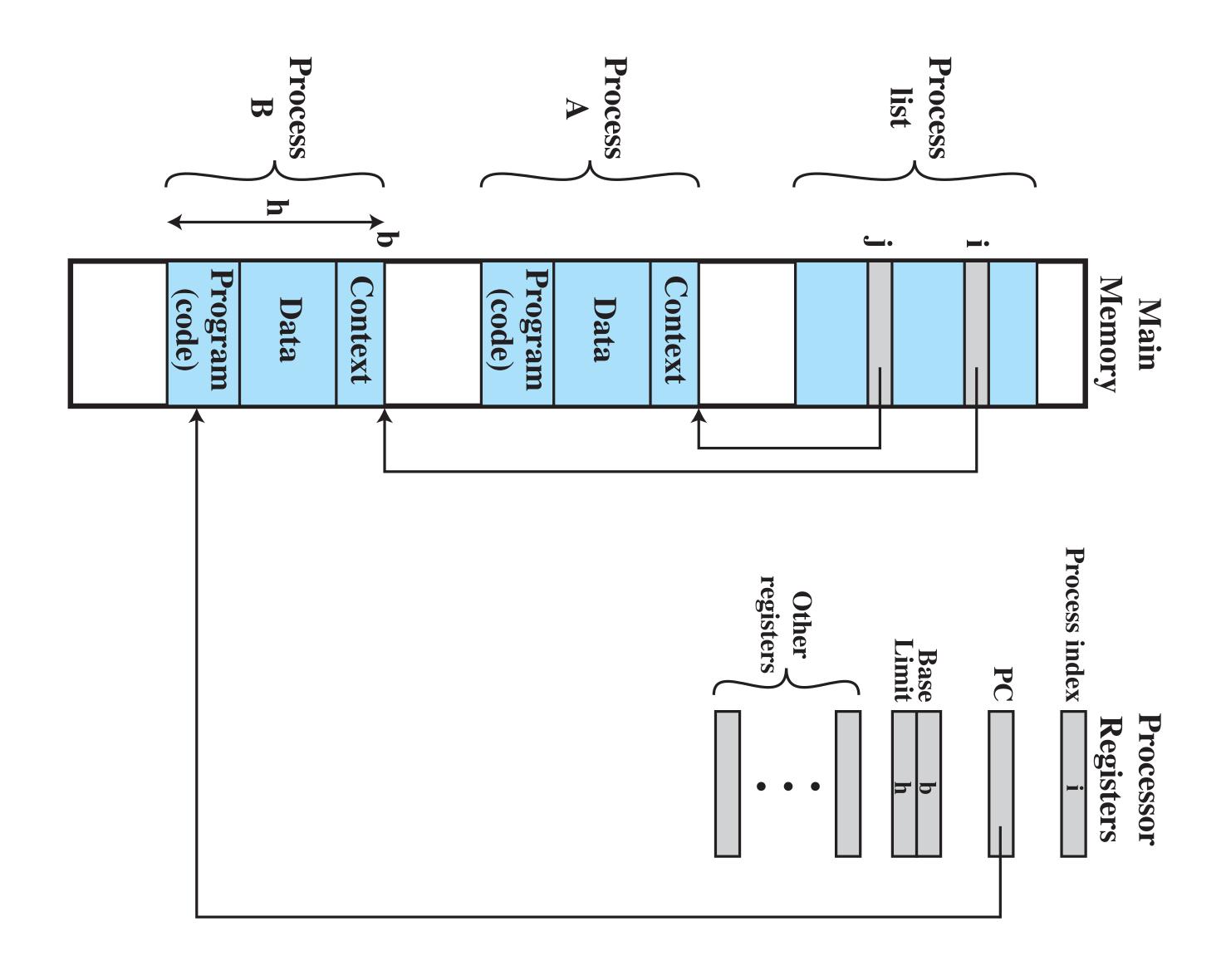


Figure 2.8 **Typical Process** Implementation

Major Achievements

- Denning et.al. [DENN80a] point out 5 major OS advances:
 - Processes
 - Memory Management
 - Information protection and security
 - Scheduling and resource management
 - System structure
- Let's look at each one...

Processes

- A program in execution
- An instance of a program running on a computer
- The entity that can be assigned to and executed on a processor
- A unit of activity characterized by a single sequential thread of execution, a current state, and an associated set of system resources

Memory Management

- Process isolation
 - non-interference between independent procs.
- Automatic allocation and management
 - should be transparent to programmer
- Support of modular programming
- Protection and access control
- Long-term storage
 after computer has been powered down

- Allows programmers to address memory from a logical point of view
- No hiatus between the execution of successive processes while one process was written out to secondary store and the successor process was read in

Virtual Memory

Virtual Memory and File System

- Implements long-term store
- Information stored in named objects called files

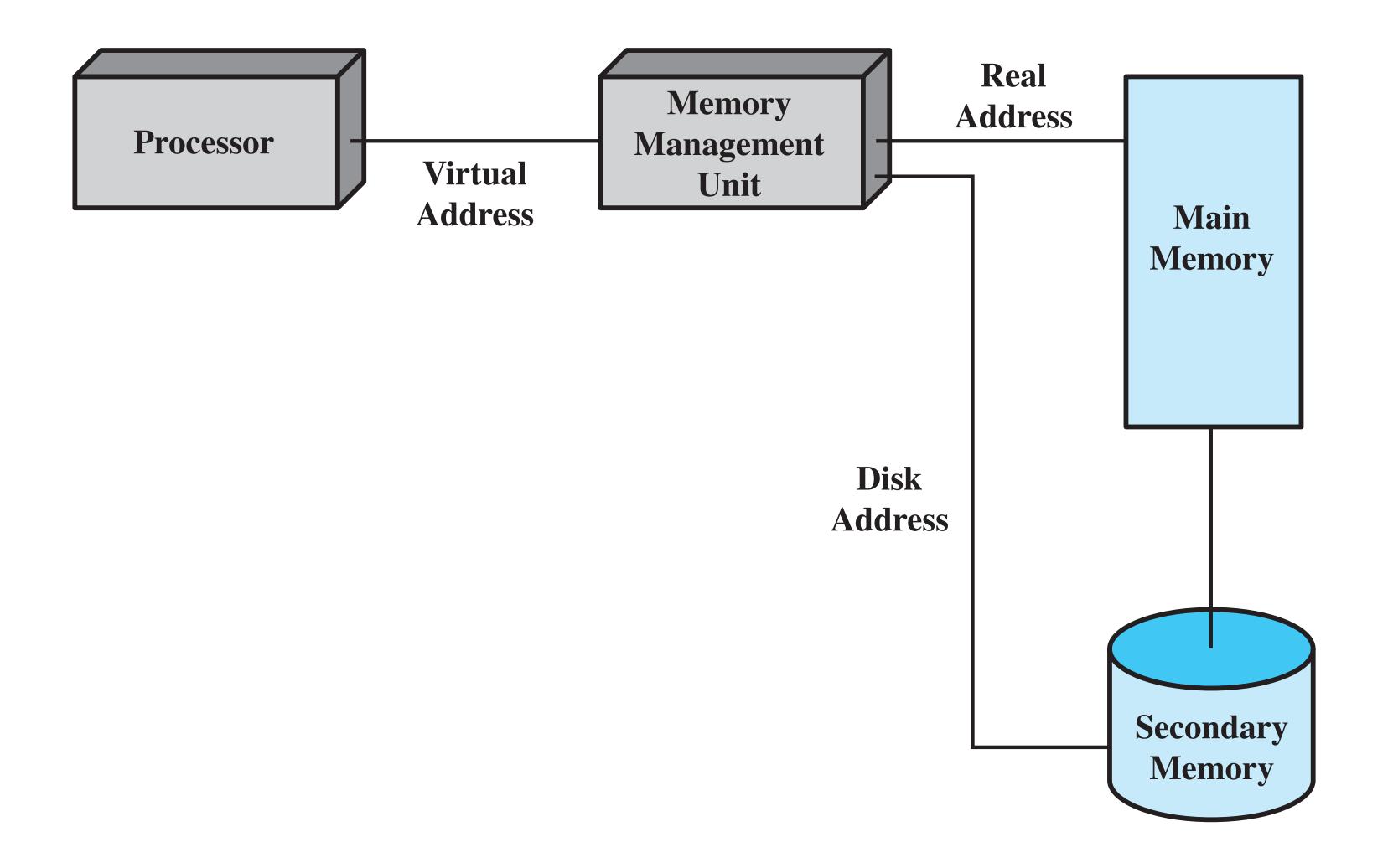


Figure 2.10 Virtual Memory Addressing

						B. 0			A.1
B.5			A.9			B.1	A.5	A.0	
B.6		A.8		A.7		B.2		A.2	
						B. 3			

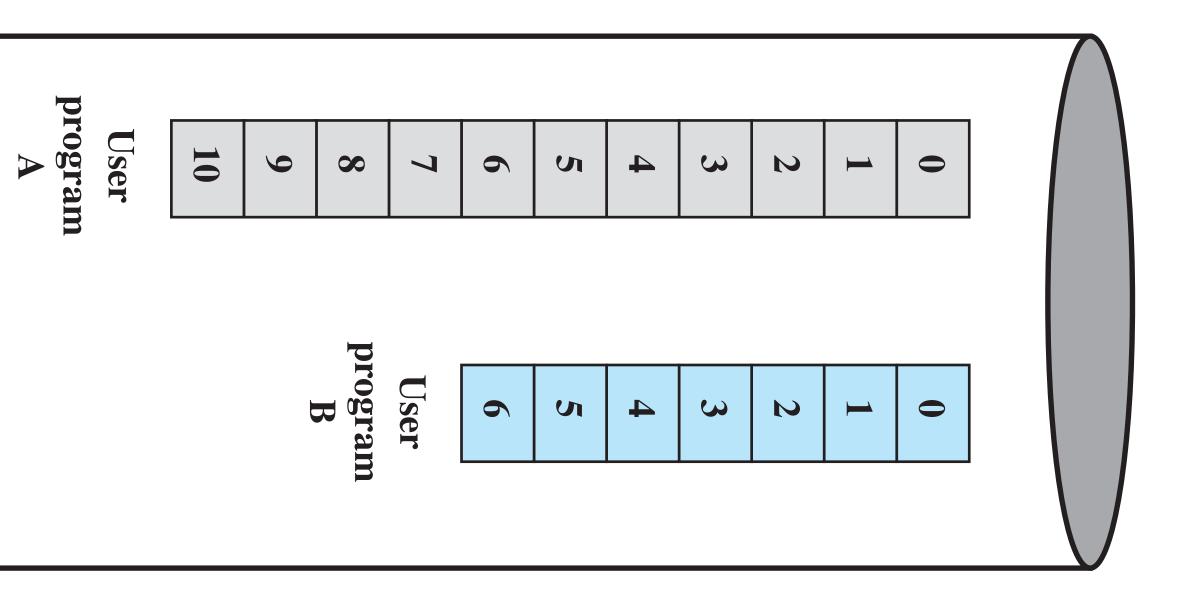
Main Memory

Main memory consists of a number of fixed-length frames, each equal to the size of a page. For a program to execute, some or all of its pages must be in main memory.

Disk

Secondary memory (disk) can hold many fixed-length pages. A user program consists of some number of pages. Pages of all programs plus the operating system are on disk, as are files.

Figure 2.9 **Virtual Memory Concepts**



Information Protection and Security

- Availability
 - Concerned with protecting the system against interruption
- Confidentiality
 - Assuring that users cannot read data for which access is unauthorized

Sequence 3

CS240

Information Protection and Security

- Data integrity
 - Protection of data from unauthorized modification
- Authenticity
 - Concerned with the proper verification of the identity of users and the validity of messages or data

Scheduling and Resource Management

- Fairness
- Give equal and fair access to resources • Differential responsiveness
 - ...but, OS also needs to discriminate among different classes of jobs
- Efficiency
 - Maximize throughput, minimize response time, and accommodate as many uses as possible

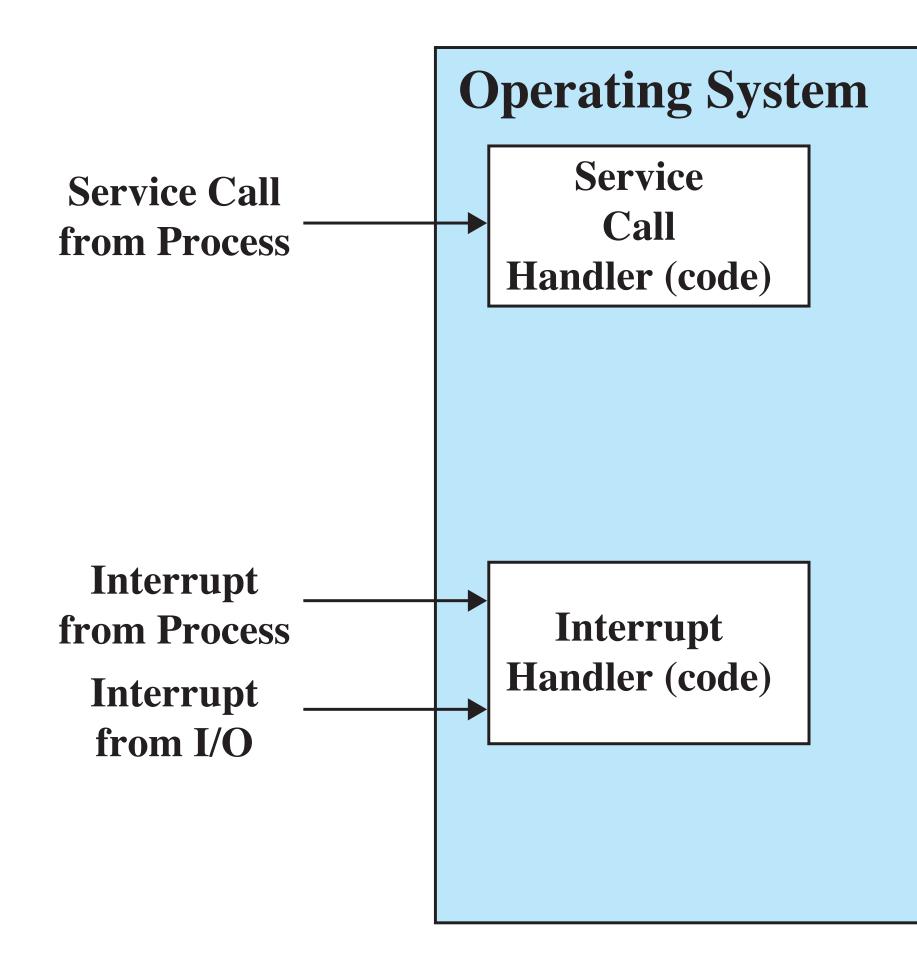
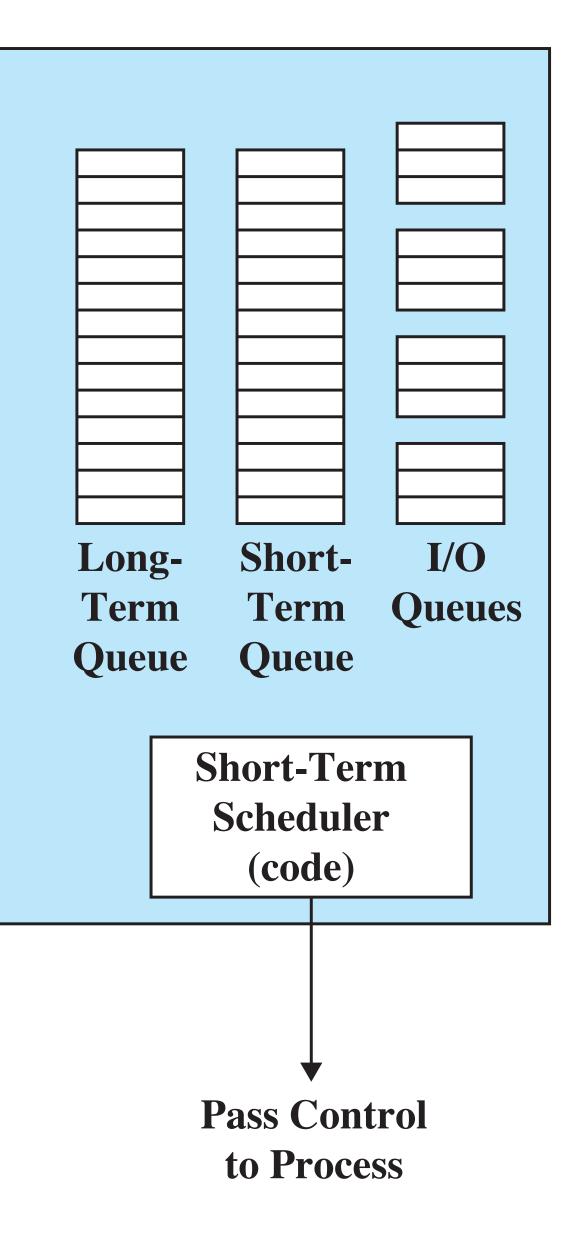
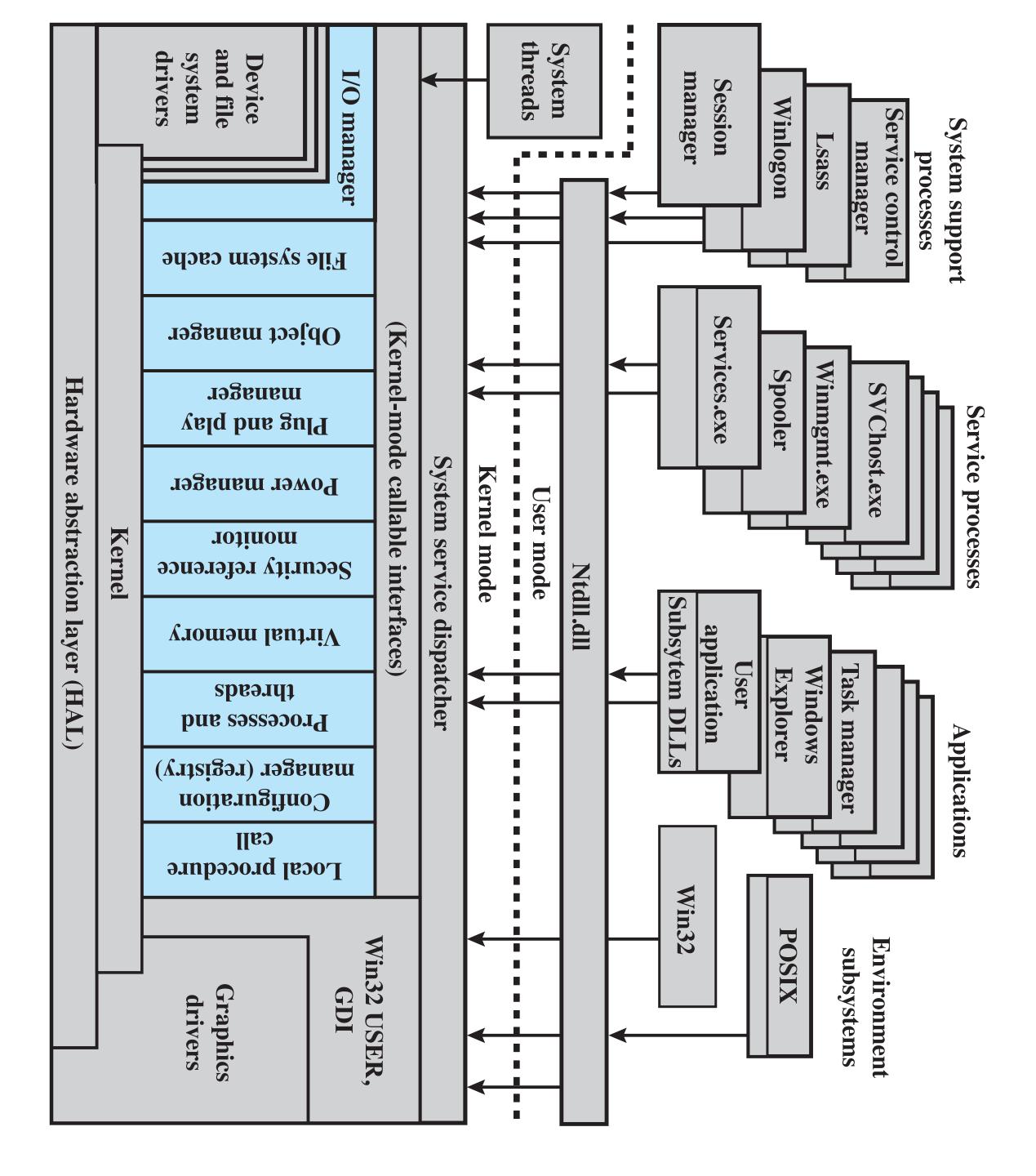


Figure 2.11 Key Elements of an Operating System for Multiprogramming





Lsass = local security authentication server POSIX = portable operating system interface GDI = graphics device interface DLL = dynamic link libraries

Colored area indicates Executive

Figure 2.14 Windows Architecture

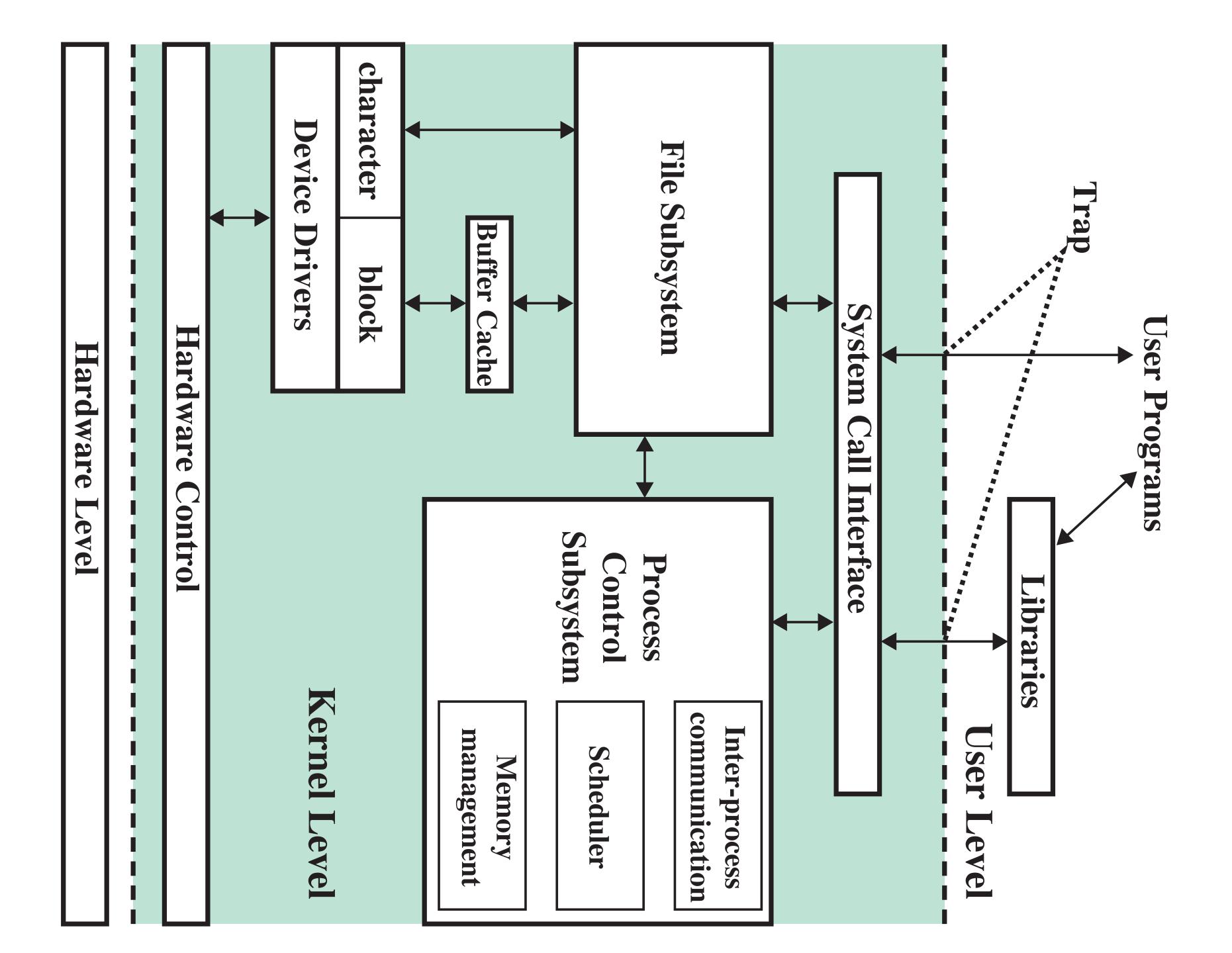


Figure 2.15 Traditional **UNIX Kernel**

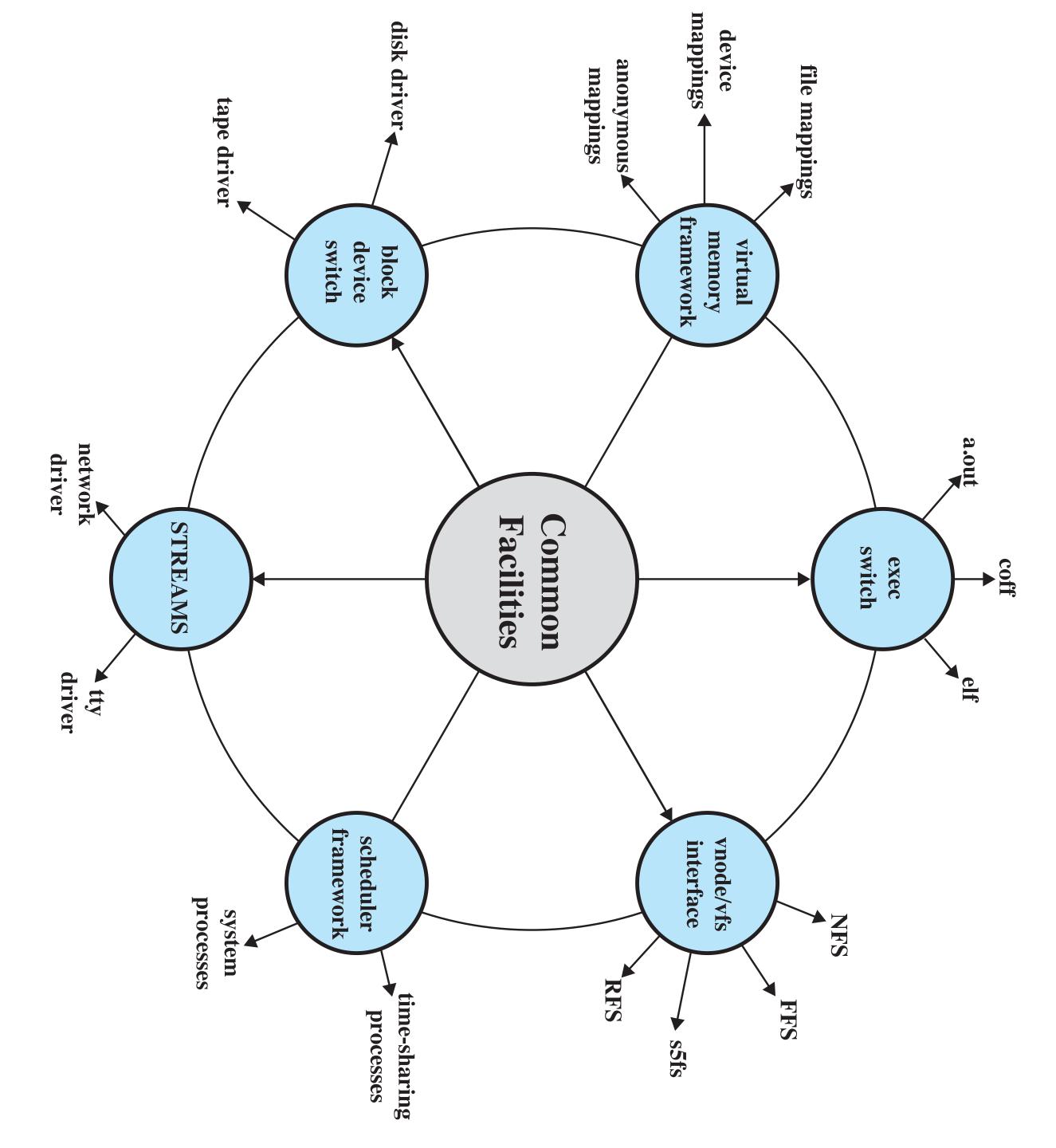


Figure 2.16 Modern UNIX Kernel

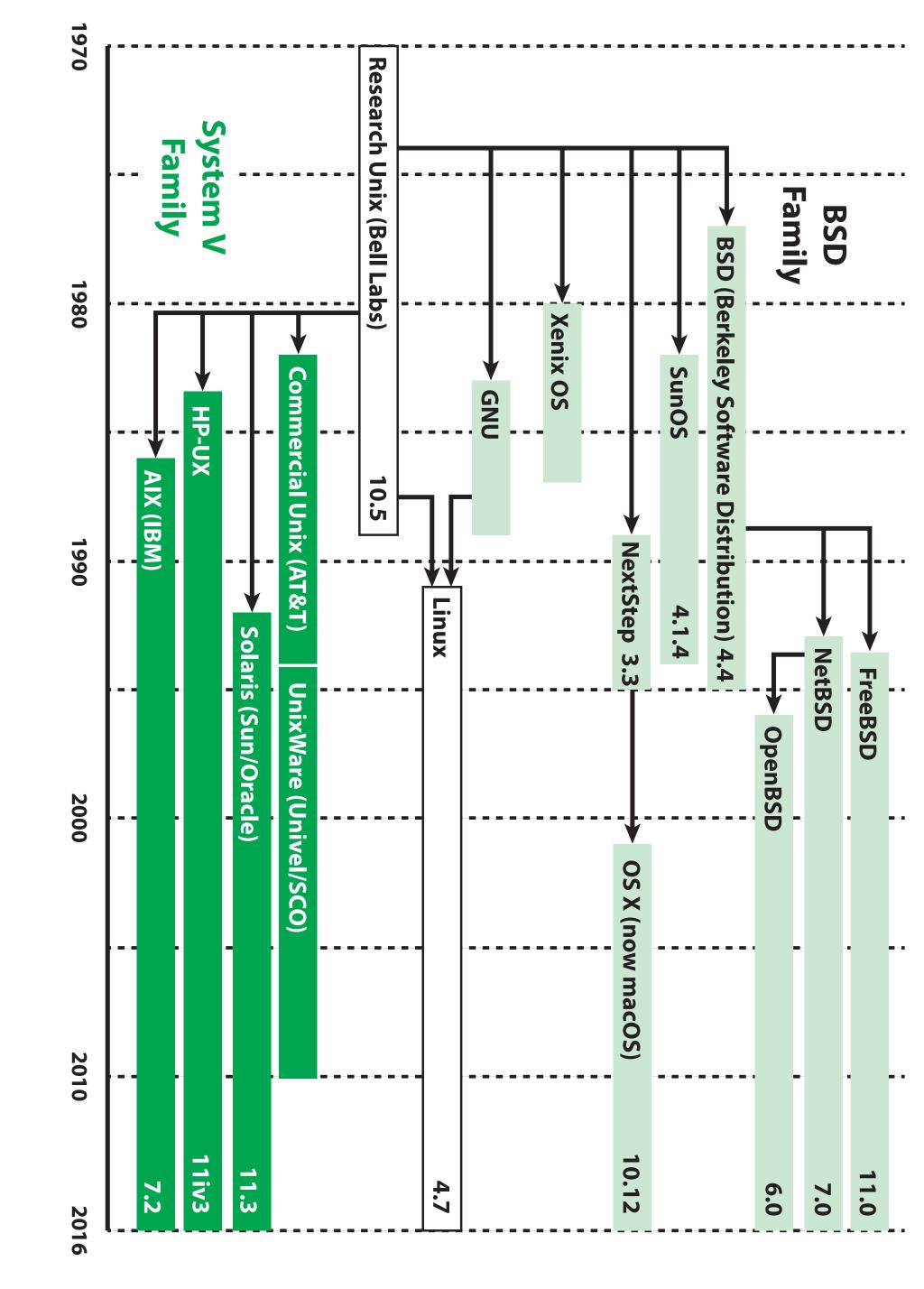
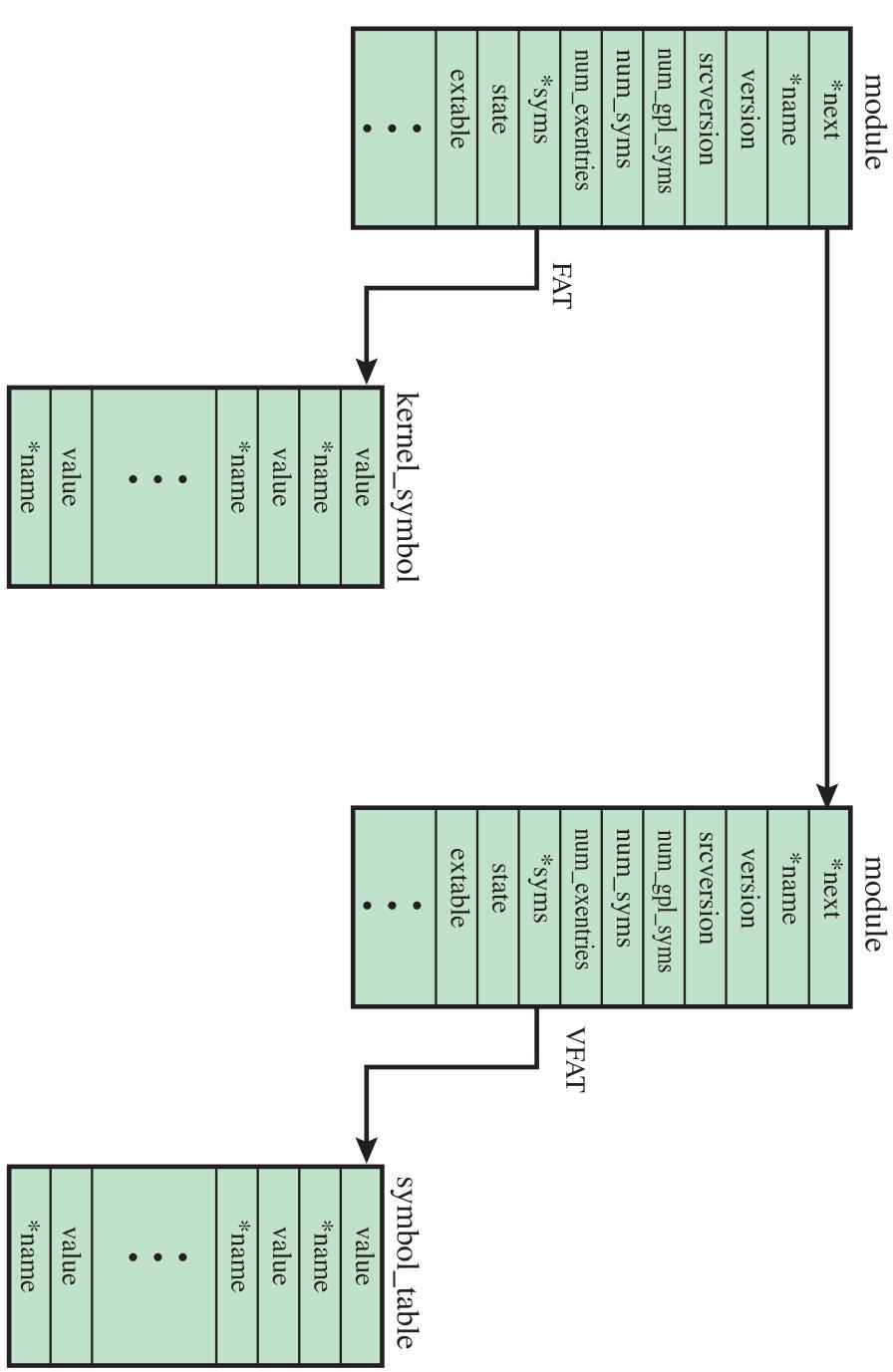


Figure 2.17 Unix Family Tree



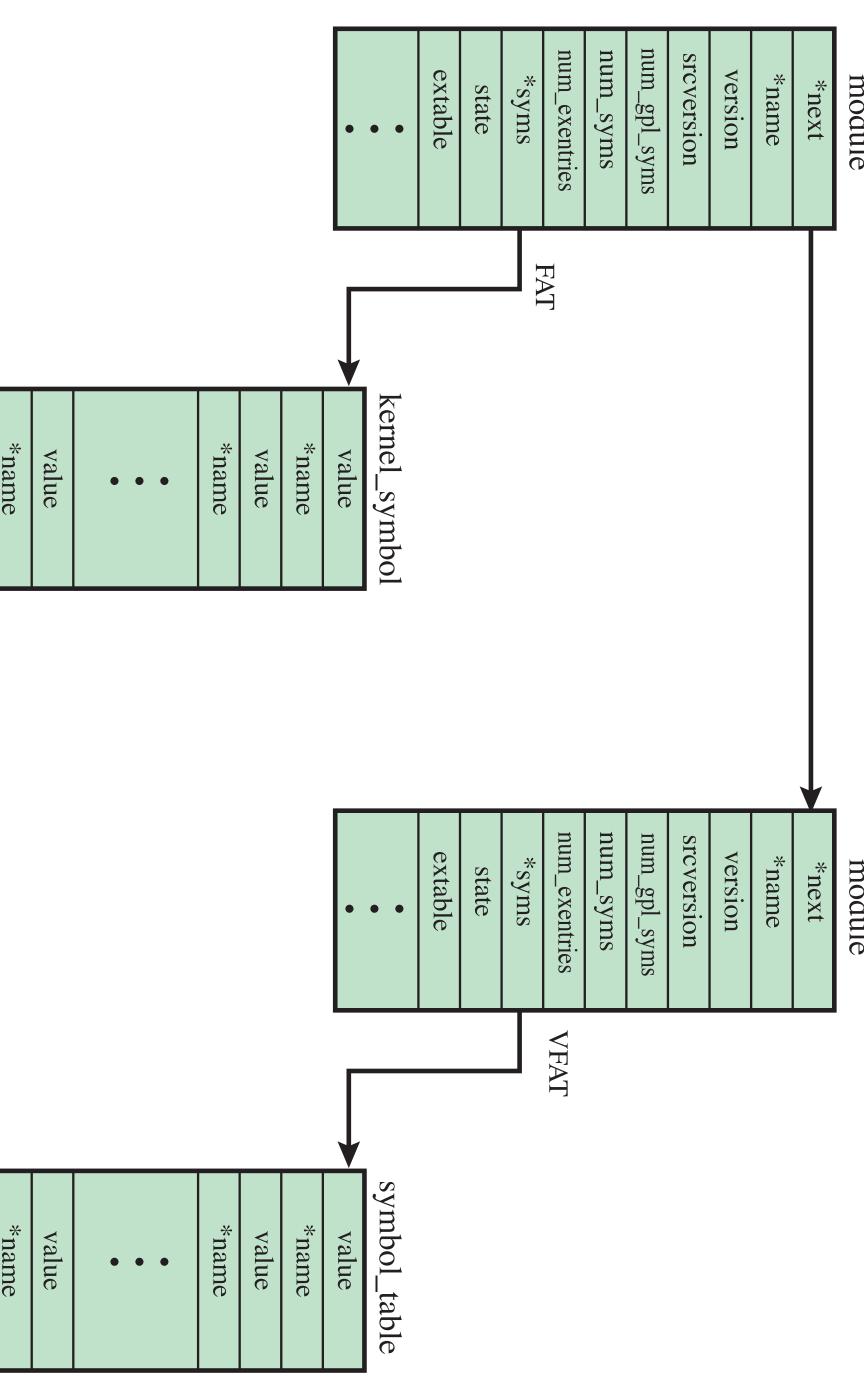


Figure 2.18 **Example List of Linux Kernel Modules**

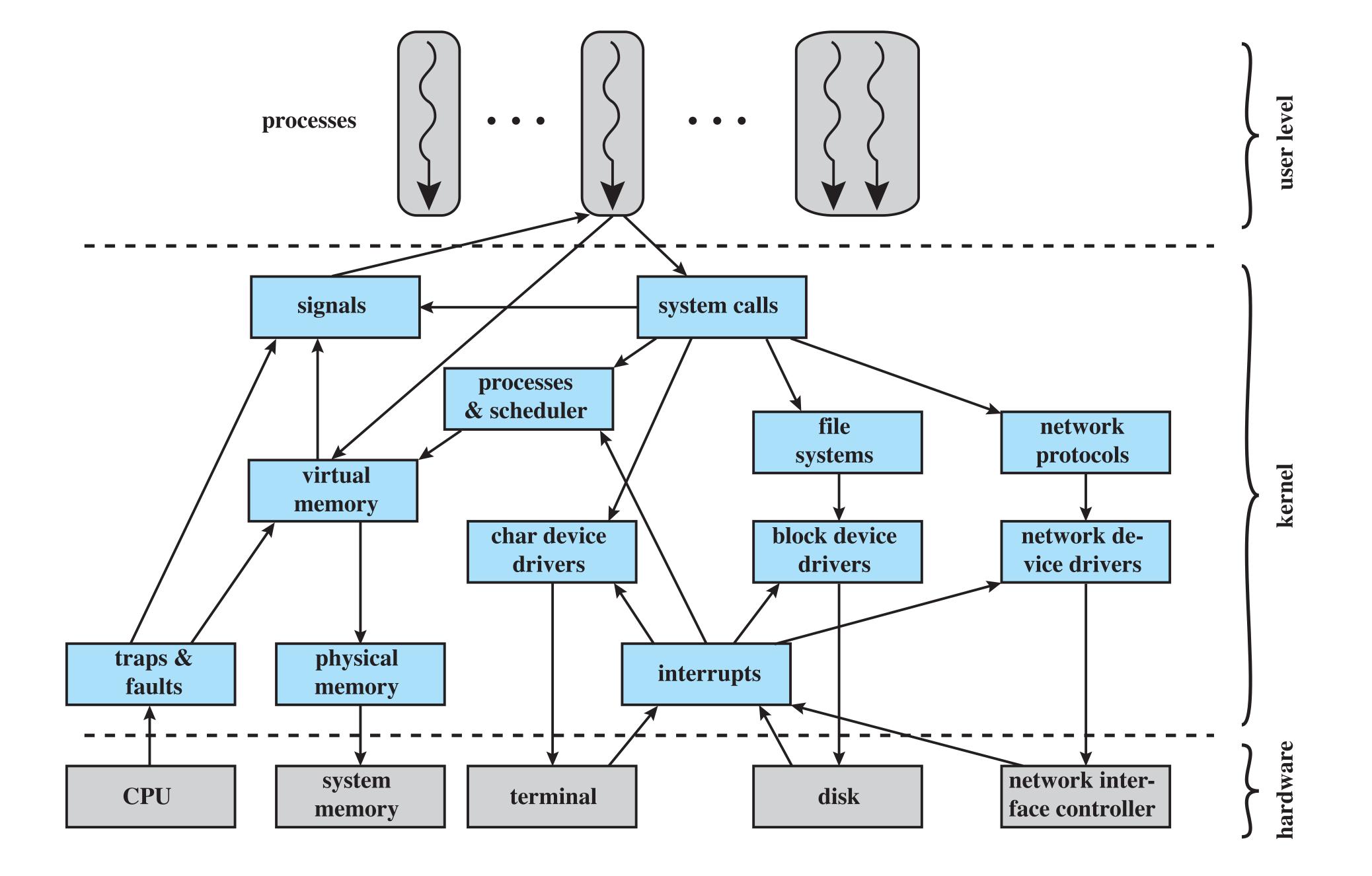
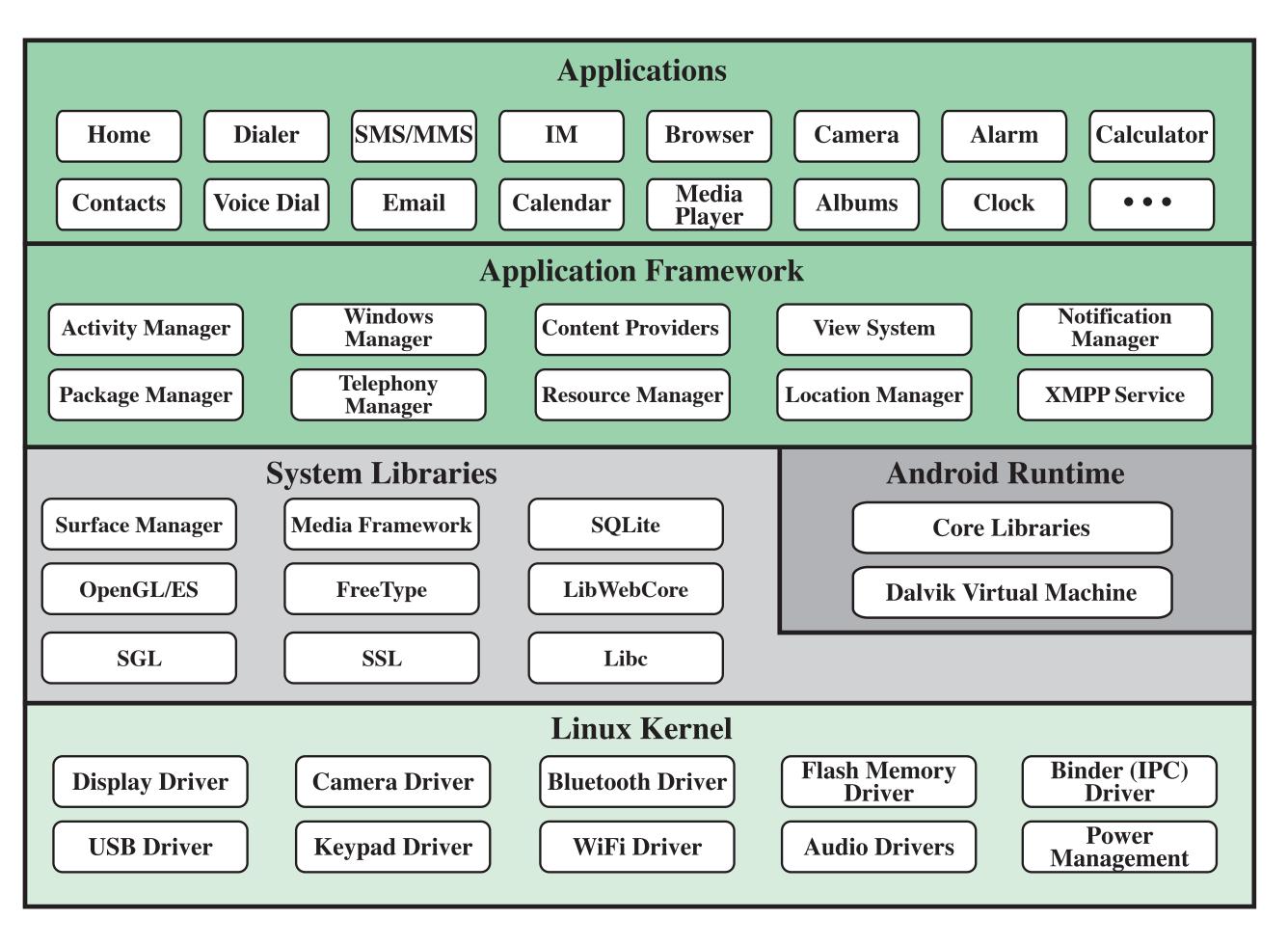


Figure 2.19 Linux Kernel Components



Implementation:



Applications, Application Framework: Java



System Libraries, Android Runtime: C and C++

Linux Kernel: C

Figure 2.20 Android Software Architecture

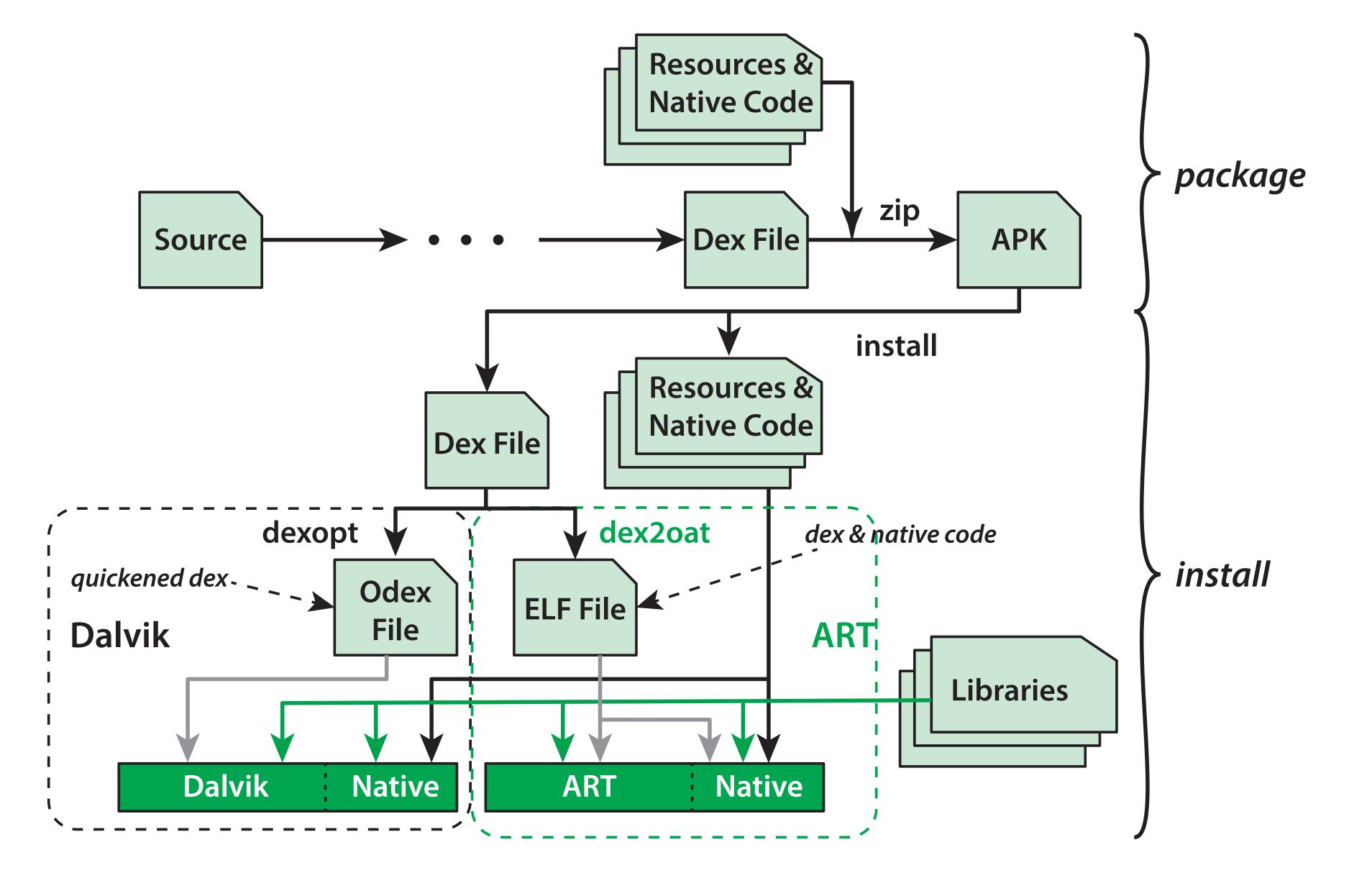


Figure 2.21 The Life Cycle of an APK

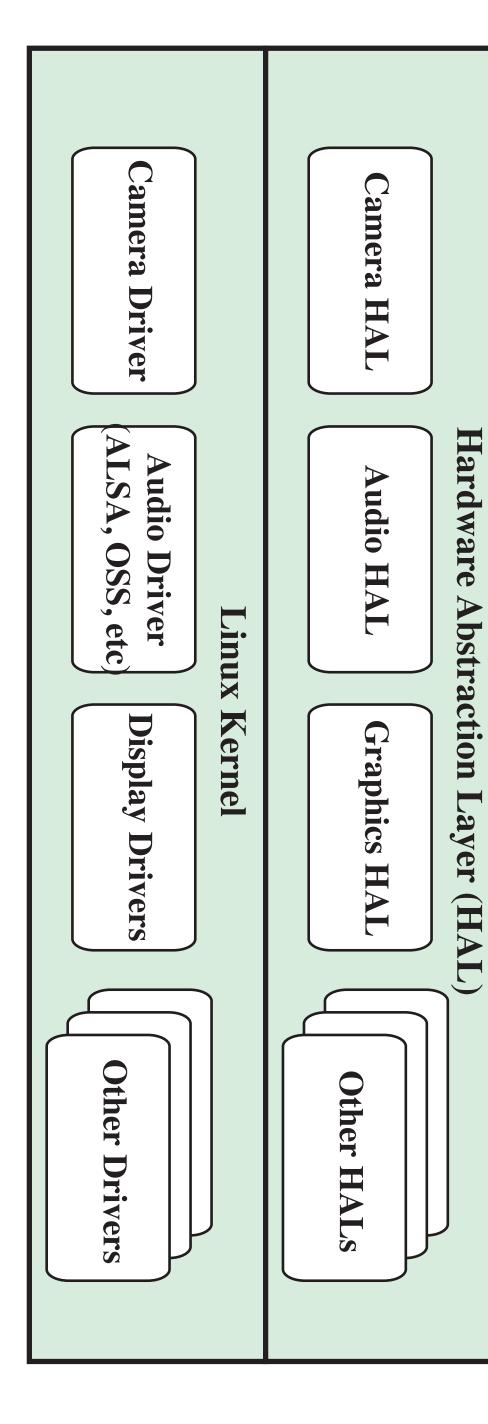


Figure 2.22 Android System Architectu Ire

AudioFlinger Camera Service Ľ Ľ **Media Server** L. L. E. 1 ł MediaPlayer **Other Media** 1 Service **Applications and Framework** Services **Android System Services** Android Runtime/Dalvik U D 9 н. t. Ľ U. Ľ. U. D Binder IPC 1 1 Manager Manager Activity L. Service Power ŧ. U, d 11 11 System Server Ľ Ц Other S Manager L Window 1 L. ervices