Deadlock

• Permanent blocking of a set of processes that either compete for system resources or communicate with each other
• No efficient solution
• Involve conflicting needs for resources by two or more processes
Figure 6.1  Illustration of Deadlock
Better illustrations 😊
Reusable Resources

- Used by only one process at a time and not depleted by that use
- Processes obtain resources that they later release for reuse by other processes
  - E.g. Processors, I/O channels, main and secondary memory, devices, and data structures such as files, databases, and semaphores
- Deadlock occurs if each process holds one resource and requests the other
Example of Deadlock

Now consider the following sequence:

\[ p_0 \ p_1 \ q_0 \ q_1 \ p_2 \ q_2 \]
Another Example of Deadlock

• Space is available for allocation of 200Kbytes, and the following sequence of events occur

<table>
<thead>
<tr>
<th>P1</th>
<th>P2</th>
</tr>
</thead>
<tbody>
<tr>
<td>⋮</td>
<td>⋮</td>
</tr>
<tr>
<td>Request 80 Kbytes;</td>
<td>Request 70 Kbytes;</td>
</tr>
<tr>
<td>⋮</td>
<td>⋮</td>
</tr>
<tr>
<td>Request 60 Kbytes;</td>
<td>Request 80 Kbytes;</td>
</tr>
</tbody>
</table>

• Deadlock occurs if both processes progress to their second request
Consumable Resources

- Created (produced) and destroyed (consumed)
- Interrupts, signals, messages, and information in I/O buffers
- Deadlock may occur if a Receive message is blocking
- May take a rare combination of events to cause deadlock
Example of Deadlock

- Deadlock occurs if receive is blocking

```
    P1
    ...  
    Receive(P2);
    ...  
    Send(P2, M1);

    P2
    ...  
    Receive(P1);
    ...  
    Send(P1, M2);
```
Resource Allocation Graphs

• Directed graph that depicts a state of the system of resources and processes

(a) Resource is requested
(b) Resource is held
Resource Allocation Graphs

(c) Circular wait

(d) No deadlock

Figure 6.5 Examples of Resource Allocation Graphs
Conditions for Deadlock

• Mutual exclusion
  – Only one process may use a resource at a time

• Hold-and-wait
  – A process may hold allocated resources while awaiting assignment of others

• No preemption
  – No resource can be forcibly removed from a process holding it
Conditions for Deadlock

• Circular wait
  – A closed chain of processes exists, such that each process holds at least one resource needed by the next process in the chain