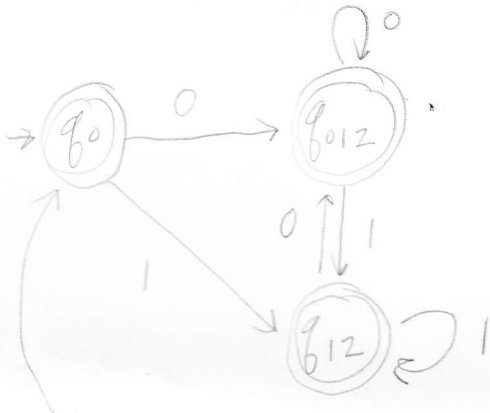
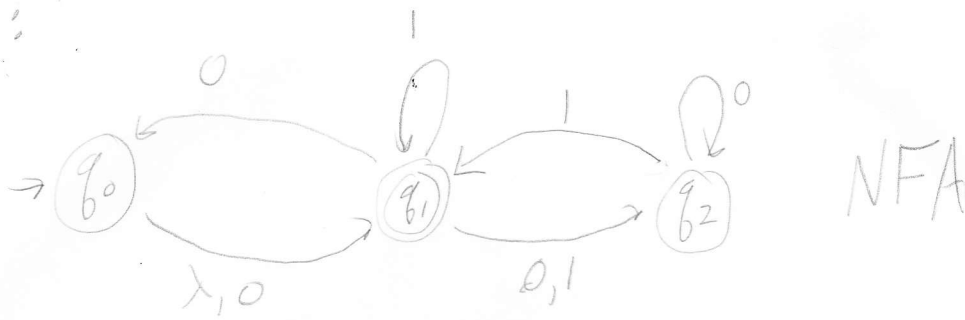


Answer Key #3

Page 62:

3)

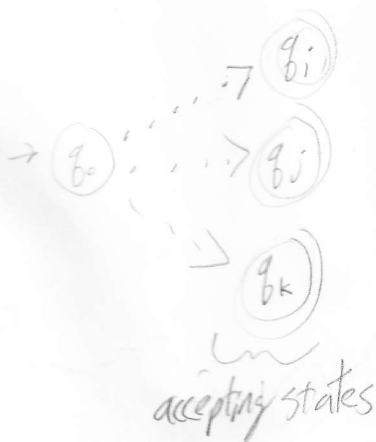


Note q_0 must be accepting because the original NFA has a λ transition to an accepting state.

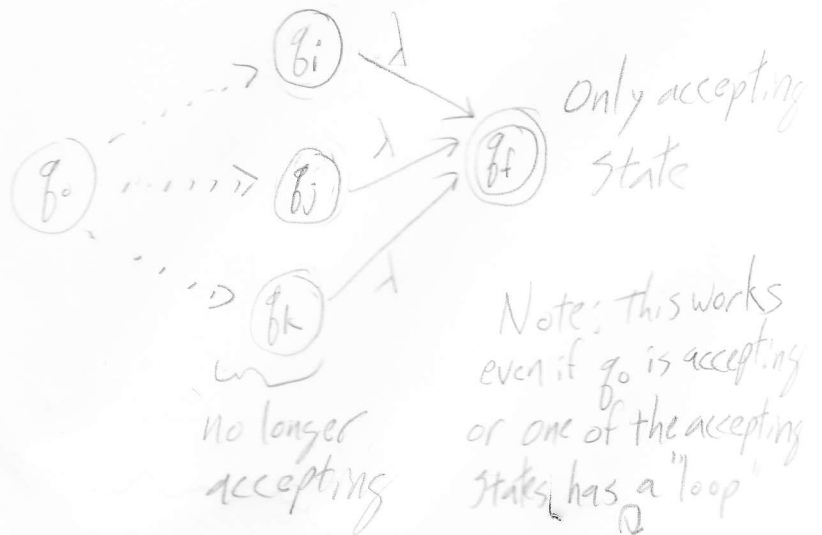
7) solved in text.

Visually it looks like:

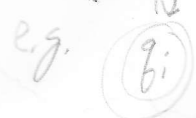
Original NFA



which becomes



Note: this works even if q_0 is accepting or one of the accepting states has a "loop"



Page 62:

11) solved in text.

Page 75: 14, 16ab

1) find all strings in $L((a+b)^*b(a+b)^*)$ of length less than 4.
recall that $(a+b)^*$ is "anything" including ϵ .

$L = \{b, ab, bb, aab, abb, bab, bbb, \\ ba, bab, baa, \\ aba, bba\}$

4) $\{a^n b^m : n \geq 3, m \text{ is even}\}$

a^n $n \geq 3$ can be done w/ 3a's plus some more: $\underbrace{aaa}_{3a's} \underbrace{a^*}_{\text{some more}}$

b^m m is even requires b's to come in pairs: $(bb)^*$

or $bb(bb)^*$ if zero is not even

So, the answer is:

$aaa^*bb(bb)^*$

also $a^*(aaa)(bb)^*bb$ or $aa(a^*)a(bb)^*b$ or...

Page 75 (cont):

$$16a) \Sigma = \{a, b, c\}$$

all strings containing exactly one a.

Note $(b+c)^*$ is any string of b's & c's (including none)

So, $(b+c)^* a (b+c)^*$

$$16b) \Sigma = \{a, b, c\}$$

all strings containing no more than 3 a's

Long version: $(b+c)^* + (b+c)^* a (b+c)^* + (b+c)^* a (b+c)^* a (b+c)^* + \dots$

no a's OR 1 a w/ any # of b's and c's around it 2 a's

Shorter:

$$(b+c)^* (\lambda + a) (b+c)^* (\lambda + a) (b+c)^* (\lambda + a) (b+c)^*$$