

Assignment #6/7/Key (4a, 4b, 4c, 5a, 14, 15a, 15b, 24, 26, Page 133: 2, 3, 7a)  
Page 122.

4a) Answer in text

$$4b) L = \{a^n b^k a^k; k \neq n+1\}$$

Assume  $N$  as in the P.L.

(Example 4.13 is a good place to look for hints)

$$\text{Choose } w = a^{N!} b a^{(N+1)!+1}$$

w/  $w = xyz$  &  $|xy| \leq N$   $y$  must be some number of  $a$ 's from the 1<sup>st</sup> part of the string, say  $y$  is  $p$   $a$ 's  
i.e.  $y = a^p$

Now to contradict the P.L. and prove  $L$  is not regular we need to show that for any  $p \leq N$  the following:

$$N! + \underbrace{(i-1)p + 1}_{\text{for the } b} = \underbrace{(N+1)! + 1}_{\text{the extra } a \text{ to balance the } b} \text{ is true for some } i$$

Solve the equation

$$N! + (i-1)p + 1 = (N+1)! + 1$$

$$(i-1)p = (N+1)! - N! \quad \text{'pull out' an } N!$$

$$(i-1)p = N!(N+1-1) = N!N$$

$$i = \frac{N!N}{p} + 1 \quad \text{which is the } i \text{ to make } w_i \notin L \text{ invalidating the P.L. and proving } L \text{ is not regular}$$

$$4c) L = \{a^n b^l a^k; n=l \text{ or } l \neq k\}$$

Note for a string not to be in  $L$  it must be the case that  $n \neq l$  and  $l = k$ , so for the proof you need to find an  $i$  that makes  $n \neq l$  and  $l = k$

Let  $N$  be as in the p.h.

Pick  $w = a^N b^N c^N$  this is in  $L$  because  $n=l$

now  $x$  must be some # of a's from beginning of the string.

$w_i$  (for any  $i > 1$ ) will have  $n \neq l$  because there are 'too many' a's and  $l = k$  because it always was.

5a) answer in text

5b)  $L = \{a^n : n \text{ is not prime}\}$  I'll be lazy and use proof-by-contradiction

Assume  $L$  is regular

Then, because the regular languages are closed under complementation  $\bar{L}$  is regular.

But by problem 5a  $\bar{L}$  is not regular, thus the assumption that  $L$  is regular must be false.

24. We know  $L_1 \cup L_2$  is regular  
and  $L_1$  is regular

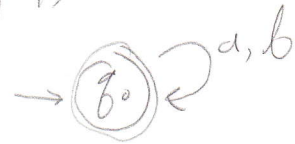
Is  $L_2$  necessarily regular? No, although it might be.

Proof:

Let  $L_1$  be all strings. Eg. over the  
alphabet  $\{a, b\}$   $L_1 = L((a+b)^*) = L(M_1)$  where

Clearly  $L_1$  is regular (I just wrote  
a regular expression and a DFA for it).

$M_1$  is



$L_1 \cup L_2 = L_1$  also regular

So, whether  $L_2$  is regular or not

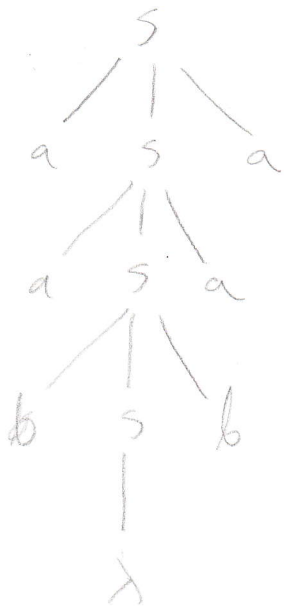
$L_1$  and  $L_1 \cup L_2$  are.

26.  $L = \{a^n b^m : n \geq 100, m \leq 50\}$

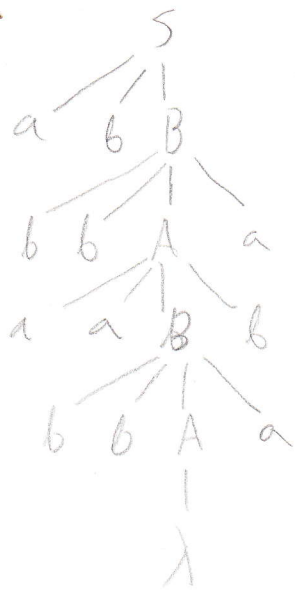
a) No. You can not use the P.L. to prove  $L$  is regular because  
the P.L. can only be used to prove that languages  
are not regular.

b) No. Because  $L$  is regular.

2



3.  $w = \underline{abbb} \underline{aabb} \underline{aba}$



Leftmost Derivation

$$\begin{aligned}
 S &\Rightarrow a b \underline{B} \Rightarrow a b b b \underline{A} a \\
 &\Rightarrow a b b b a a \underline{B} b a \\
 &\Rightarrow a b b b a a b b \underline{A} a b a \\
 &\Rightarrow a b b b a a b b a b a
 \end{aligned}$$

Note that there is only one variable per string so leftmost (& rightmost) is automatic.

7a) in text