

1 Yacc

A Yacc specification describes a context free grammar (CFG), that can be used to generate a parser.

Elements of a CFG:

1. Terminals: tokens and literal characters,
2. Variables (nonterminals): syntactical elements,
3. Production rules, and
4. Start rule.

Lexical Analyzer scans the input stream and converts sequences of characters into tokens. Identifies tokens in input string (stream).

Format of a production rule:

```
symbol: definition
      {action}
      ;
```

Example: $A \rightarrow Bc$ is written in yacc as `a: b 'c';`

1.1 yacc File Format

```
definitions

%%

grammar rules (productions and associated actions)

%%

user routines (C functions)
```

1.1.1 definitions

```
%{
C code
%}

%definition(s)
```

1.1.2 rules

```
rule action
rule action
rule action
```

Declarations—define tokens and their characteristics

`%token`: declare names of tokens

`%prec`: assign precedence to a rule

`%left`: define left-associative operators

`%right`: define right-associative operators

`%nonassoc`: define operators that may not associate with themselves

`%type`: declare the type of variables

`%union`: declare multiple data types for semantic values

`%start`: declare the start symbol (default is the first variable in rules)

`%{`

C declarations directly copied to the resulting C program

`%}` (E.g., variables, types, macros)

Although right-recursive rules can be used in yacc, left-recursive rules are preferred, and, in general, generate more efficient parsers.

Invoking yacc on a file, e.g.,

```
yacc yaccFile.y
```

generates two output files:

y.tab.c and y.tab.h

An application can be built using:

```
cc -o app lex.yy.c y.tab.c -ll -ly
```

1.2 Conflicts

Pointer model: A pointer moves (right) on the RHS of a rule while input tokens and variables are processed.

```
%token A B C
%%
start: A B C /* after reading A: start: A B C */
```

When all elements on the right-hand side are processed (pointer reaches the end of a rule), the rule is *reduced*. The pointer then returns to the rule it was called.

Conflict: There is a conflict if a rule is reduced when there is more than one pointer.

yytext is an external variable that contains the matched string.

1.3 Simple Calculator

1.3.1 calc1.1

```
%{
/* calc1.1
*/

#include <stdio.h>

#include "y.tab.h"

extern int yylval;
%}

%%
```

```

[0-9]+    {
            yylval = atoi(yytext);
            return INTEGER;
        }

[-+\n]    return *yytext;

[ \t]     { ; /* skip whitespace */ }

.         yyerror( "invalid character" );

%%

int yywrap( void )
{
    return 1;
}

```

1.3.2 calc1.y

```

%{
#include <stdio.h>
%}

%token INTEGER

%%

program:
    program expr '\n' { printf( "= %d\n", $2); }
    |
    ;

expr:
    INTEGER           { $$ = $1; }
    | expr '+' expr   { $$ = $1 + $3; }
    | expr '-' expr   { $$ = $1 - $3; }
    ;

```

```
%%  
  
int yyerror( char *s )  
{  
    fprintf( stderr, "%s\n", s );  
    return 0;  
}  
  
int main()  
{  
    yyparse();  
    return 0;  
}
```

\$\$ left hand side value
\$n nth argument on right hand side

1.3.3 Building the Calculator

```
# buildCalc1.sh  
#  
# Create simple calculator  
#  
# make y.tab.c and y.tab.h (-d option creates header)  
yacc -d calc1.y  
# make lex.yy.c  
lex calc1.l  
# compile and link C files  
cc -o calc1 y.tab.c lex.yy.c -ly -ll  
  
% ./buildCalc1.sh  
yacc: 4 shift/reduce conflicts.
```

1.3.4 Running the Calculator

```
% ./calc1  
1 + 2  
= 3
```

```
2 *3
invalid character
syntax error
```

References

- [Johnson79] Stephen C. Johnson, “Yacc: Yet Another Compiler-Compiler,” *UNIX Programmer’s Manual*, volume 2, Holt, Rinehart, and Winston, pages 353–387, 1979.
- [LM&B90] Levine, Mason & Brown, *lex & yacc*, O’Reilly & Associates, 1990.
- [S&FB85] Schreiner and Friedman, *Introduction to Compiler Construction with Unix*, Prentice-Hall, 1985.
- [LnYP03] The Lex & Yacc Page, <http://dinosaur.compilertools.net/>