Part of this discussion is based on the paper

- Petri Nets: Properties, Analysis and Applications
- by Tadoa Murata, Proc. IEEE, Vol. 77, No. 4, April 1989.

Petri Nets

- graphical and mathematical modeling tool
- tool for describing systems characterized as being:
 - » concurrent, asynchronous, distributed, parallel, nondeterministic and/or stochastic

History

- 1962: Carl Adam Petri's submitted his dissertation at the Uni.
 Darmstadt, Germany
- 1970: early development was published by A.W. Host and in the records of the 1970 Project MAC Conference on Concurrent Systems and Parallel Computation
- 1970-75: Computation Structure Group and MIT was most active
- 1975: conference on Petri Nets and Related Methods at MIT
- **1979:** 135 researchers assembled in Hamburg, Germany, for 2-week advanced course on General Net Theory of Processes and Systems
- **1980:** first European Workshop on Applications and Theory of Petri Nets, Strasbourg, France.
- check out Murata's paper for the extensive literature discussion

General:

- directed, weighted, bipartite graph
- two kinds of notes (Places P, Transitions T)
- arcs from P to T or from T to P
- arcs have integer weights
- non-negative Place weights are called tokens

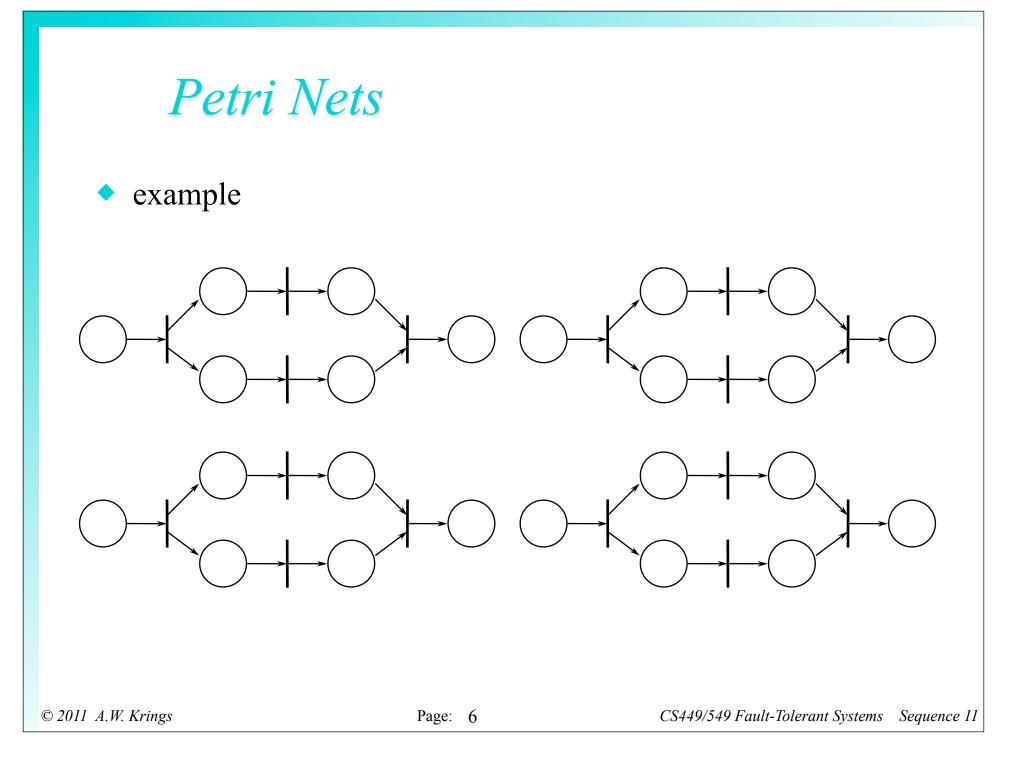
- A Petri Net is a 5-touple PN={P,T,A,W,M0}
- Place Set $P = \{p_1, p_2, ..., p_m\}$
 - finite set of places
 - condition = place
 - one condition or set of atomic conditions
 - symbol (

• Transition Set $T = \{t_1, t_2, ..., t_n\}$

- finite set of transitions
- action = transition
- one action or set of atomic transitions
- symbol

• Arc Set $A \subseteq (P \times T) \cup (T \times P)$

- set of directed arcs
- edge of graph = arc
- symbol —
- Weight Function $W = A \rightarrow \{1, 2, 3, ...\}$
 - weights are associated with arcs
- Initial Marking $M_0 = P \rightarrow \{0, 1, 2, ...\}$
 - the initial assignment of tokens to places



Dynamic Behavior

- during simulation of a petri net the state of the net may change
- change of state:
 - » transitions can be enabled
 - » enabled transitions may fire
 - » firing transition changes the marking of the net
 - » the marking is the "snap-shot" of all the tokens

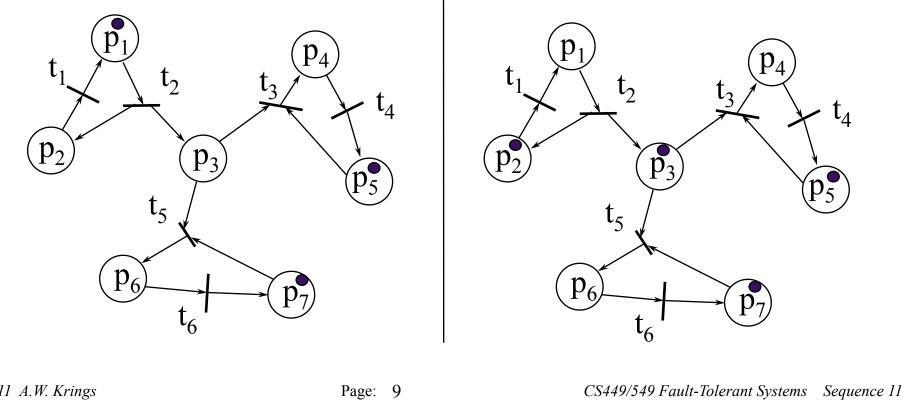
Firing rules

- A transition *T* is said to be *enabled* if each input place *P* is marked with at least W(P,T) tokens
 - » W(P,T) is the weight of the arc from P to T
- An enabled transition may or may not fire (depending on whether or not the event actually takes place).
- A *firing* of an enabled transition T removes W(P,T) tokens from each input place P of T, and adds W(T,P) tokens to each output place P of T
 - » W(T,P) is the weight of the arc from T to P
- Common misconception: When a transition fires, it does not move tokens
 - » i.e. the number of tokens in the system is not necessarily constant

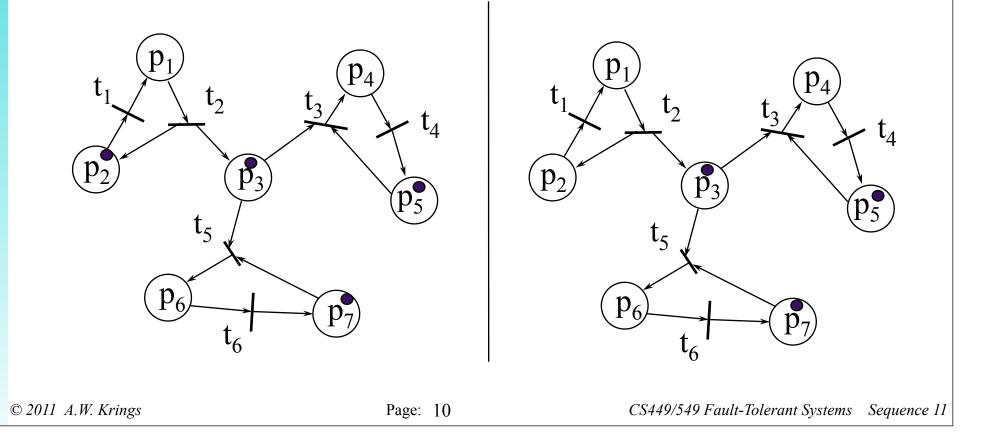


• Example: assume the following initial marking

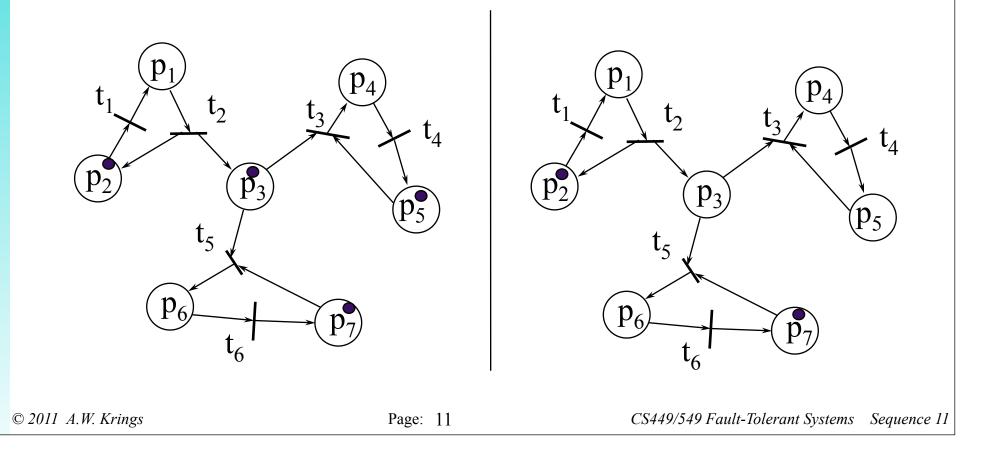
Only one transition is enabled, i.e. t_2 _



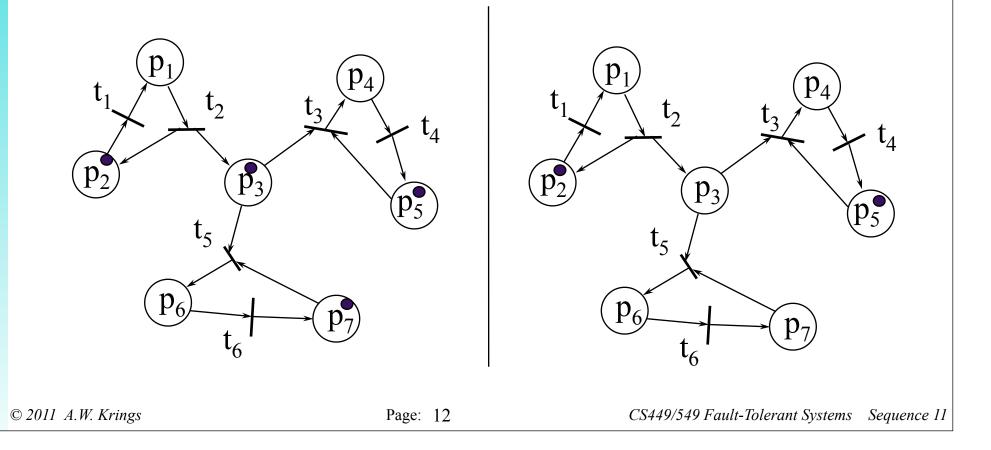
- Now several transitions are enabled, i.e. $t_1 t_3$ and t_5
- if t_1 fires first



- if t_3 fires first

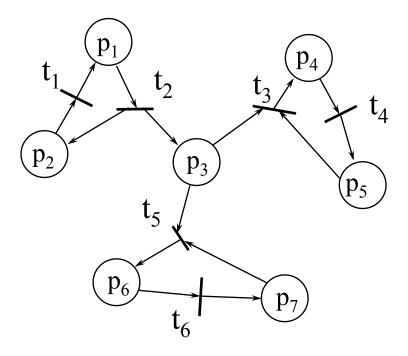


- if t_5 fires first
- t_3 and t_5 are said to be in conflict





- what could this Petri net represent?



Marking: Number and placement of tokens

- let $m_i = \#$ of tokens in place p_i
- then marking

 $M = \{m_1, m_2, ..., m_n\}$

- marking -- system state
- Advantage: economy of model
 - » e.g. assume net with 6 places
 - we limit each place to maximal 1 token
 - then there are 2⁶ possible markings
 - \blacksquare => 64 states
 - thus Petri Nets are a lot smaller than state diagrams, i.e. Markov chains

Firing rules

- transition 1,3 and 4 are enabled

