## The Internet Protocol

## Chapter 14 in Stallings $10^{\text {th }}$ Edition

```
Communication Network
    A facility that provides a data transfer service among devices attached to the
network.
Internet
    A collection of communication networks interconnected by bridges and/or routers.
Intranet
    An internet used by a single organization that provides the key Internet applications,
especially the World Wide Web. An intranet operates within the organization for internal
purposes and can exist as an isolated, self-contained internet, or may have links to the
Internet
```


## Subnetwork

```
Refers to a constituent network of an internet. This avoids ambiguity because the entire internet, from a user's point of view, is a single network.
```


## End System (ES)

```
A device attached to one of the networks of an internet that is used to support enduser applications or services
Intermediate System (IS)
A device used to connect two networks and permit communication between end systems attached to different networks
Bridge
An IS used to connect two LANs that use similar LAN protocols. The bridge acts as an address filter, picking up packets from one LAN that are intended for a destination on another LAN and passing those packets on. The bridge does not modify the contents of the packets and does not add anything to the packet. The bridge operates at layer 2 of the the packets
OSI model.
```


## Router

```
An IS used to connect two networks that may or may not be similar. The router employs an internet protocol present in each router and each end system of the network. The router operates at layer 3 of the OSI model. Page 2
```



Figure 14.1 TCP/IP Concepts


## Connectionless Operation

- Internetworking involves connectionless operation at the level of the Internet Protocol (IP)

- Initially developed for the DARPA internet project
- Protocol is needed to access a particular network


## Connectionless <br> Internetworking

- Advantages
—Flexibility
—Robust
—No unnecessary overhead
- Unreliable
-Not guaranteed delivery
-Not guaranteed order of delivery
- Packets can take different routes
—Reliability is responsibility of next layer up (e.g. TCP)



## Connectionless Internetworking

- Connectionless internet facility is flexible
- IP provides a connectionless service between end systems
-Advantages:
- Is flexible
- Can be made robust
- Does not impose unnecessary overhead


## IP Design Issues

- Routing
- Datagram lifetime
- Fragmentation and reassembly
- Error control
- Flow control


# The Internet as a Network 


(a) Packet-switching network architecture
$\mathrm{S}=$ station
$\mathrm{P}=$ packet switching node
$\mathrm{R}=$ router
$\mathrm{T}=$ transmission link
$\mathrm{N}=$ network

(b) Internetwork architecture

## Routing

- Routing table indicates next router to which datagram is sent
- Can be static or dynamic

ES / routers maintain routing tables

Source routing

- Source specifies route to be followed
- Can be useful for security and priority
- Each router appends its internet address to a list of addresses in the datagram
- Useful for testing and debugging purposes

Route recording

## Datagram Lifetime

- If dynamic or alternate routing is used the potential exists for a datagram to loop indefinitely
-Consumes resources
-Transport protocol may need upper bound on lifetime of a datagram
- Can mark datagram with lifetime
- When lifetime expires, datagram is discarded


## Fragmentation and Re-assembly

- Protocol exchanges data between two entities
- Lower-level protocols may need to break data up into smaller blocks, called fragmentation
- Reasons for fragmentation:
- Network only accepts blocks of a certain size
- More efficient error control and smaller retransmission units
- Fairer access to shared facilities
- Smaller buffers
- Disadvantages:
- Smaller buffers
- More interrupts and processing time


## Fragmentation and Reassembly



## IP Fragmentation

- IP reassembles at destination only
- Uses fields in header


Figure 14.4 Fragmentation Example

## Error and Flow Control

Error control
-Discarded datagram identification is needed
-Reasons for discarded datagrams include:

- Lifetime expiration
- Congestion
- FCS error

Flow control
-Allows routers to limit the rate they receive data
-Send flow control packets requesting reduced data flow

## Internet Protocol (IP) v4

- Defined in RFC 791
- Part of TCP/IP suite
- Two parts

Specification of interface with a higher layer

Specification of actual protocol format and mechanisms

## IP Services

- Primitives
-Specifies functions to be performed
-Form of primitive implementation dependent
-Send-request transmission of data unit
-Deliver-notify user of arrival of data unit
- Parameters
—Used to pass data and control information


## IP Parameters

- Source and destination addresses
- Protocol
- Type of Service
- Identification
- Don't fragment indicator
- Time to live
- Data length
- Option data
- User data


## IP Options





Figure 14.6 IPv4 Address Formats


## Subnets and Subnet Masks

- Allows arbitrary complexity of internetworked LANs within organization
- Insulate overall internet from growth of network numbers and routing complexity
- Site looks to rest of internet like single network
- Each LAN assigned subnet number
- Host portion of address partitioned into subnet number and host number
- Local routers route within subnetted network
- Subnet mask indicates which bits are subnet number and which are host number

Table 14.2
IPv4 Addresses and Subnet Masks

|  | Binary Representation | Dotted Decimal |
| :--- | :--- | :--- |
| IP address | 11000000.11100100 .00010001 .00111001 | 192.228 .17 .57 |
| Subnet mask | 11111111.11111111 .11111111 .11100000 | 255.255 .255 .224 |
| Bitwise AND of address <br> and mask (resultant <br> network/subnet number) | 11000000.11100100 .00010001 .00100000 | 192.228 .17 .32 |
| Subnet number | 11000000.11100100 .00010001 .001 | 1 |
| Host number | 00000000.00000000 .00000000 .00011001 | 25 |

(a) Dotted decimal and binary representations of IPv4 address and subnet masks

|  | Binary Representation | Dotted Decimal |
| :--- | :--- | :--- |
| Class A default mask | 11111111.00000000 .00000000 .00000000 | 255.0 .0 .0 |
| Example Class A mask | 11111111.11000000 .00000000 .00000000 | 255.192 .0 .0 |
| Class B default mask | 11111111.11111111 .00000000 .00000000 | 255.255 .0 .0 |
| Example Class B mask | 11111111.11111111 .11111000 .00000000 | 255.255 .248 .0 |
| Class C default mask | 11111111.11111111 .11111111 .00000000 | 255.255 .255 .0 |
| Example Class C mask | 11111111.11111111 .11111111 .11111100 | 255.255 .255 .252 |

(b) Default subnet masks

CS420/520 Axel Krings
Page 26
Sequence 14

## Inter-domain Routing

- Classless Inter-Domain Routing - CIDR
-Introduced in 1993 by the Internet Engineering Task Force
-Goal was to slow the growth of routing tables on routers across the Internet, and to help slow the rapid exhaustion of IPv4 addresses
—CIDR appends a "/" character to the address and the decimal number of leading bits of the routing prefix -Example:
- 192.168.1.0/24 for IPv4
- 2001:db8::/32 for IPv6


Figure 14.7 Example of Subnetworking

## Internet Control Message Protocol (ICMP)

- RFC 792
- Provides a means for transferring messages from routers and other hosts to a host
- Provides feedback about problems
- Datagram cannot reach its destination
- Router does not have buffer capacity to forward
- Router can send traffic on a shorter route
- Encapsulated in IP datagram
-Hence not reliable

(a) Destination Unreachable; Time Exceeded; Source Quench

(b) Parameter Problem

(c) Redirect

(d) Echo, Echo Reply

(e) Timestamp

(g) Address Mask Request

(h) Address Mask Reply

Figure 14.8 ICMP Message Formats

## Common ICMP Messages

- Destination unreachable
- Time exceeded
- Parameter problem
- Source quench
- Redirect
- Echo and echo reply
- Timestamp and timestamp reply
- Address mask request and reply


## Address Resolution Protocol (ARP)

Need MAC address to send to LAN host

## Manual

Included in network address
Use central directory
Use address resolution protocol

ARP (RFC 826) provides dynamic IP to Ethernet address mapping

Source broadcasts ARP request
Destination replies with ARP response

## IP Next Generation

Address space
exhaustion:

- Two level addressing
(network and host) wastes
space
- Network addresses used even
if not connected
- Growth of networks and the

Internet

- Extended use of TCP/IP
- Single address per host


## IPv6 RFCs

- RFC 1752 - Recommendations for the IP Next Generation Protocol
-Requirements
-PDU formats
-Addressing, routing security issues
- RFC 2460 - overall specification
- RFC 4291 - addressing structure


## IPv6 Enhancements

- Expanded 128 bit address space
- Improved option mechanism
-Most not examined by intermediate routes
- Dynamic address assignment
- Increased addressing flexibility
-Anycast and multicast
- Support for resource allocation
-Labeled packet flows




## IPv6 Flow Label

- Related sequence of packets
- Special handling
- Identified by source and destination address plus flow label
- Router treats flow as sharing attributes
- May treat flows differently
- Alternative to including all information in every header
- Have requirements on flow label processing


## IPv6 Addresses

- 128 bits long
- Assigned to interface
- Single interface may have multiple unicast addresses


## Three types of addresses:

- Unicast - single interface address
- Anycast - one of a set of interface addresses
- Multicast - all of a set of interfaces

Table 14.3
IPv6 Address Space Usage

| Address Type | Binary Prefix | IPv6 Notation | Fraction of address <br> space |
| :--- | :--- | :--- | :--- |
| Embedded IPv4 <br> address | $00 \ldots 111111111111$ <br> 1111 <br> $(96$ bits $)$ | $\because:$ FFFF/96 | $2^{-96}$ |
| Loopback | $00 \ldots 1$ <br> $(128$ bits $)$ | $\because: 1 / 128$ | $2^{-128}$ |
| Link-local unicast | 1111111010 | FE80::/10 | $1 / 1024$ |
| Multicast | 11111111 | FF00::/8 | $2 / 256$ |
| Global unicast | Everything else |  |  |

## Hop-by-Hop Options

- Must be examined by every router
-If unknown discard/forward handling is specified
- Next header
- Header extension length
- Options
-Pad1
—PadN
—Jumbo payload
—Router alert

(b) Fragmentation header

Figure 14.10 IPv6 Extension Headers

## Fragmentation Header

- Fragmentation only allowed at source
- No fragmentation at intermediate routers
- Node must perform path discovery to find smallest MTU of intermediate networks
- Set source fragments to match MTU
- Otherwise limit to 1280 octets


## Routing Header

- Contains a list of one or more intermediate nodes to be visited on the way to a packet's destination



## Destination Options Header



## Virtual Private Network (VPN)

- Set of computers interconnected using an unsecure network
- e.g. linking corporate LANs over Internet
- Using encryption and special protocols to provide security
- Eavesdropping
- Entry point for unauthorized users
- Proprietary solutions are problematical
- Development of IPSec standard



## Benefits of IPsec

- Provides strong security for external traffic
- Resistant to bypass
- Below transport layer hence transparent to applications
- Can be transparent to end users
- Can provide security for individual users if needed


## IPsec Functions

Authentication header (AH)

- For authentication only

Encapsulating Security Payload (ESP)

- For combined authentication/encryption

A key exchange function

- Manual or automated


## VPNs usually need combined function

## Summary

- Principles of internetworking
- Requirements
- Connectionless operation
- Internet protocol operation
- Operation of a connectionless internetworking scheme
- Design issues
- Internet protocol
—IP services
- Internet protocol
— IP addresses
- ICMP
- ARP

CS420/520 Axel Krings

- IPv6
- Structure
- Header
- Addresses
- IP next generation
- Hop-by-hop options header
- Fragment header
- Routing header
-Destination options header
- VPNs and IP security
- IPsec
- Applications of IPsec
- Benefits of IPsec
- IPsec functions

